

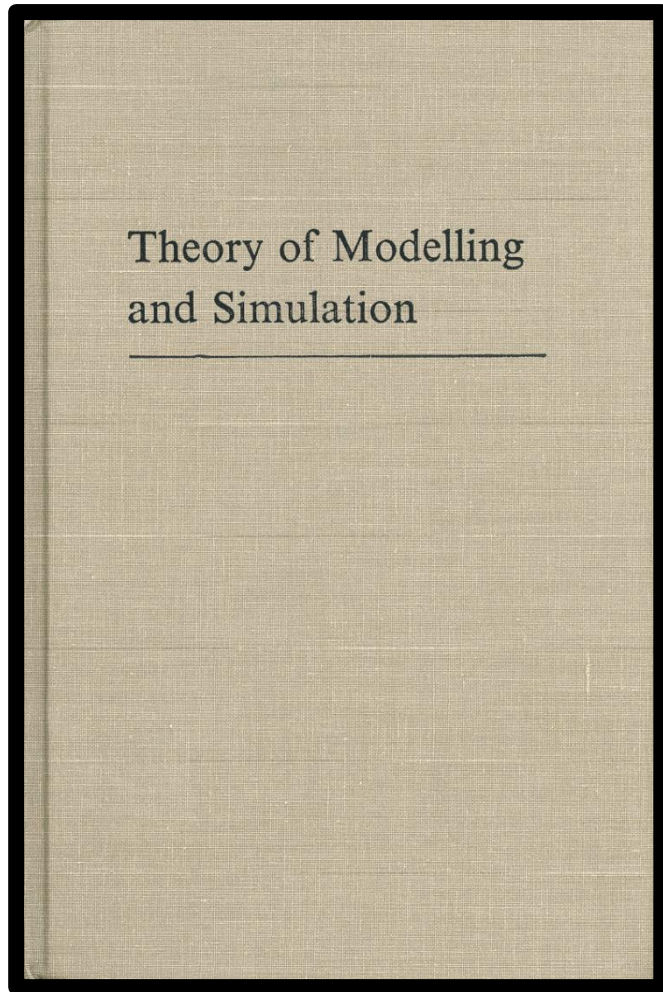
Classic DEVS

Practical Modelling Techniques Using PythonPDEVS

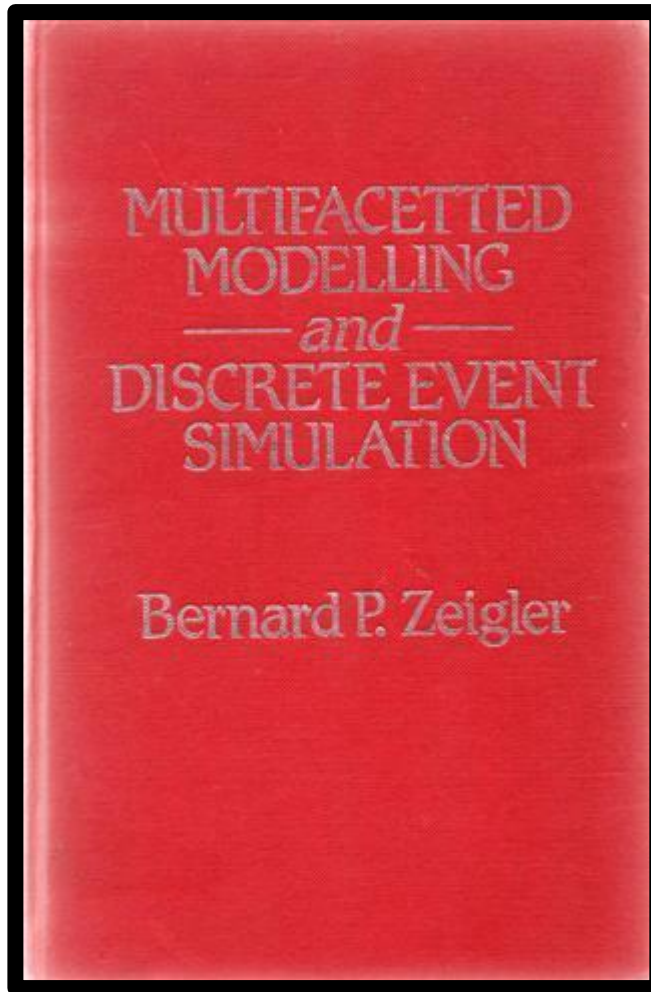
Yentl Van Tendeloo, Hans Vangheluwe

Introduction

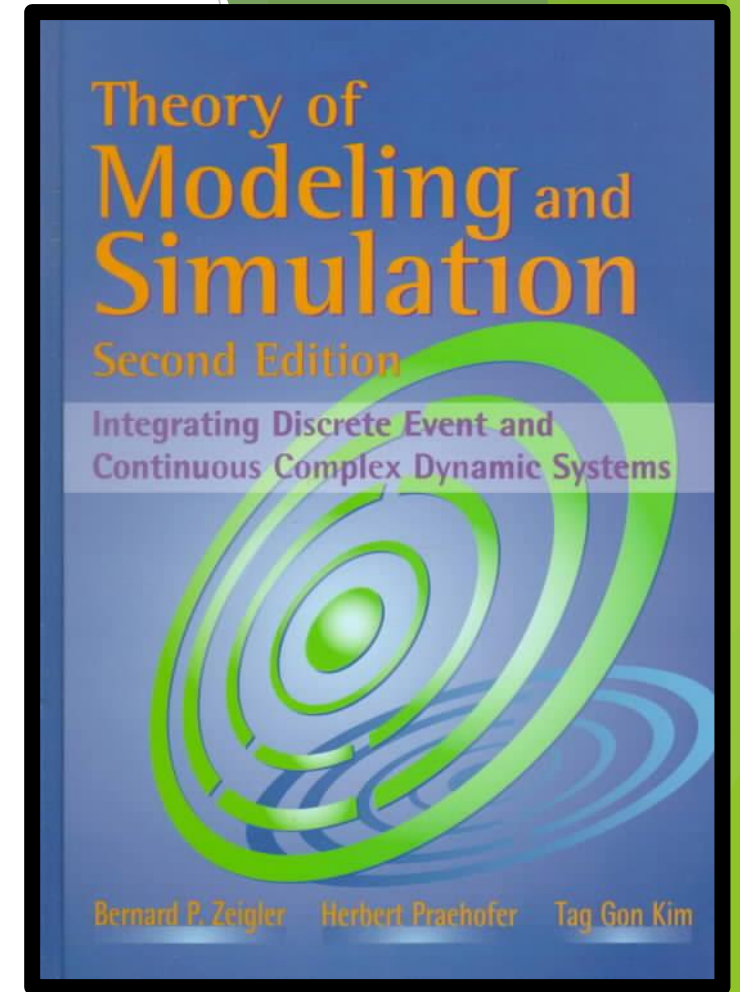
Brief recap of introductory DEVS tutorial



Bernard P. Zeigler.
Theory Of Modelling And Simulation.
1st ed. Wiley, 1976.



Bernard P. Zeigler.
*Multifaceted Modelling and
Discrete Event Simulation.*
1st ed. Academic Press, 1984.



Bernard P. Zeigler, Herbert Praehofer,
and Tag Gon Kim.
Theory Of Modelling And Simulation.
2nd ed. Academic Press, 2000.

An overview of PythonPDEVS

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Yentl Van Tendeloo and Hans Vangheluwe.
An Overview of PythonPDEVS.
In Proceedings of Journées DEVS
Francophones (JDF), pages 59-66, 2016.

Methodology

An evaluation of DEVS simulation tools

Yentl Van Tendeloo^{1*} and Hans Vangheluwe^{1,2,3*}

Abstract

DEVS is a popular formalism for modeling complex dynamic systems using a discrete-event abstraction. Owing to its popularity, and the simplicity of the simulation kernel, a number of tools have been constructed by academia and industry. However, each of these tools has distinct design goals and a specific programming language implementation. Consequently, each supports a specific set of formalisms, combined with a specific set of features. Performance differs significantly between different tools. We provide an overview of the current state of eight different DEVS simulation tools: ADEVS, CD++, DEVS-Suite, MS4 Me, PowerDEVS, PythonPDEVS, VLE, and X-S-Y. We compare supported formalisms, compliance, features, and performance. This paper aims to help modelers in deciding which tools to use to solve their specific problems. It further aims to help tool builders, by showing the aspects of their tools that could be extended in future tool versions.

Simulation



Simulation: Transactions of the Society for
Modeling and Simulation International
1-19

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Yentl Van Tendeloo and Hans Vangheluwe.
An Evaluation of DEVS Simulation Tools.
Simulation: Transactions of the Society
for Modeling and Simulation International.
2017, 93(2): 103-121



Our presentation uses initialized DEVS models, which contain an initial total state. This was left implicit in the original DEVS specification.

Extending the DEVS Formalism with Initialization Information

Yentl Van Tendeloo* Hans Vangheluwe*†
{Yentl.VanTendeloo,Hans.Vangheluwe}@uantwerpen.be

DEVS is a popular formalism to model system behaviour using a discrete-event abstraction. The main advantages of DEVS are its rigorous and precise specification, as well as its support for modular, hierarchical construction of models. DEVS frequently serves as a simulation “assembly language” to which models in other formalisms are translated, either giving meaning to new (domain-specific) languages, or reproducing semantics of existing languages. Despite this rigorous definition of its syntax and semantics, initialization of DEVS models is left unspecified in both the Classic and Parallel DEVS formalism definition. In this paper, we extend the DEVS formalism by including an initial total state. Extensions to syntax as well as denotational (closure under coupling) and operational semantics (abstract simulator) are presented. The extension is applicable to both main variants of the DEVS formalism. Our extension is such that it adds to, but does not alter the original specification. All changes are illustrated by means of a traffic light example.

Keywords: Classic DEVS, Parallel DEVS, Experimentation, Initialization

Yentl Van Tendeloo and Hans Vangheluwe. [*Extending the DEVS Formalism with Initialization Information*](#), 2018. ArXiv e-prints.

Atomic DEVS Specification

$$M = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

X : set of input events

Y : set of output events

S : set of sequential states

$q_{init} : Q$

$$Q = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$$

$\delta_{int} : S \rightarrow S$

$\delta_{ext} : Q \times X \rightarrow S$

$\lambda : S \rightarrow Y \cup \{\phi\}$

$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$

Atomic DEVS Specification

$$M = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

X : set of input events

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$q_{init} : Q$

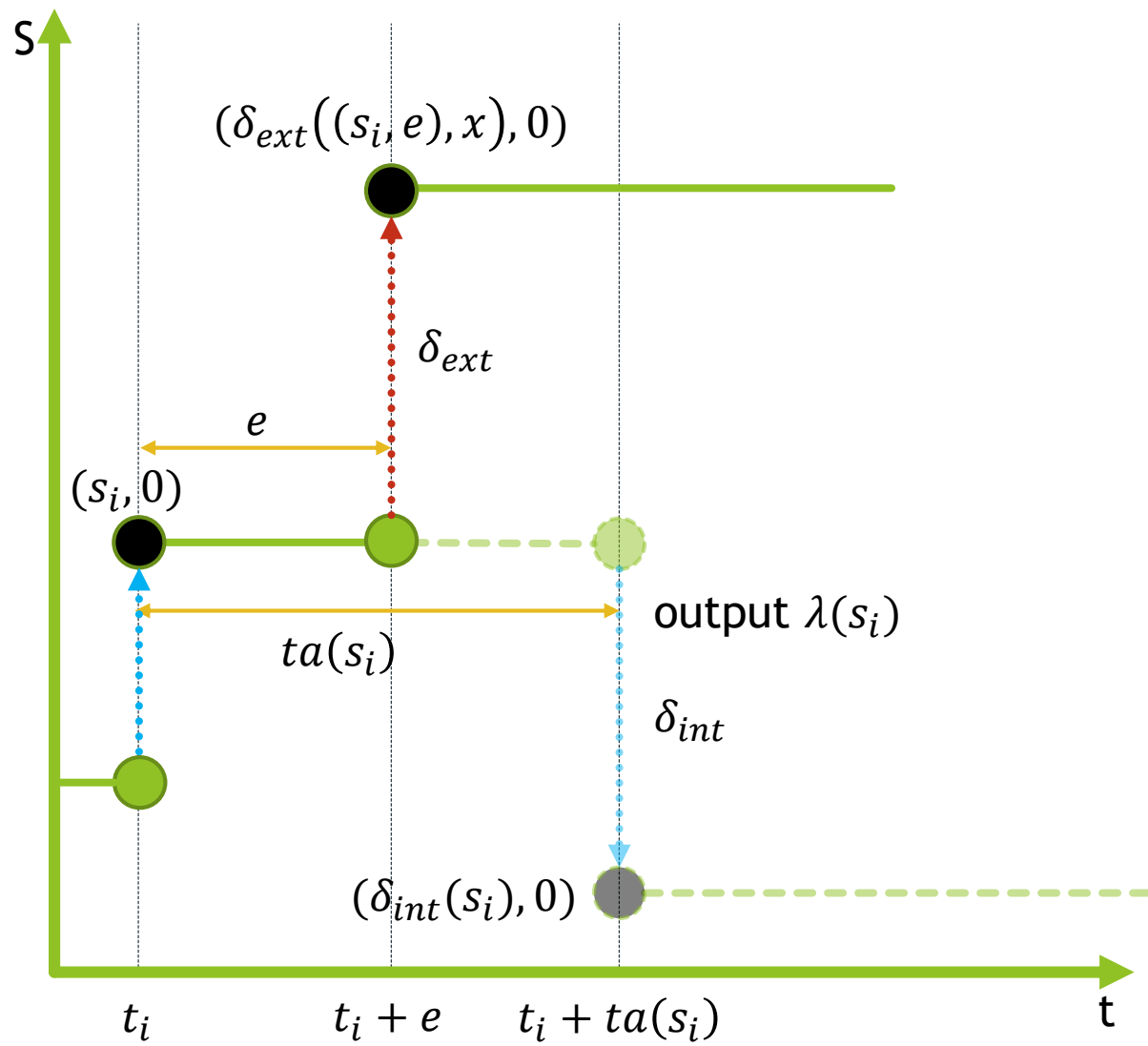
$$Q = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$$

$\delta_{int} : S \rightarrow S$

$\delta_{ext} : Q \times X \rightarrow S$

$\lambda : S \rightarrow Y \cup \{\phi\}$

$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$



Coupled DEVS Specification

$$C = \langle X_{self}, Y_{self}, D, MS, IS, ZS, select \rangle$$

$$MS = \{M_i | i \in D\}$$

$$M_i = \langle X_i, Y_i, S_i, q_{init,i}, \delta_{int,i}, \delta_{ext,i}, \lambda_i, ta_i \rangle, \forall i \in D$$

$$IS = \{I_i | i \in D \cup \{self\}\}$$

$$\forall i \in D \cup \{self\} : I_i \subseteq D \cup \{self\}$$

$$\forall i \in D \cup \{self\} : i \notin I_i$$

$$ZS = \{Z_{i,j} | i \in D \cup \{self\}, j \in I_i\}$$

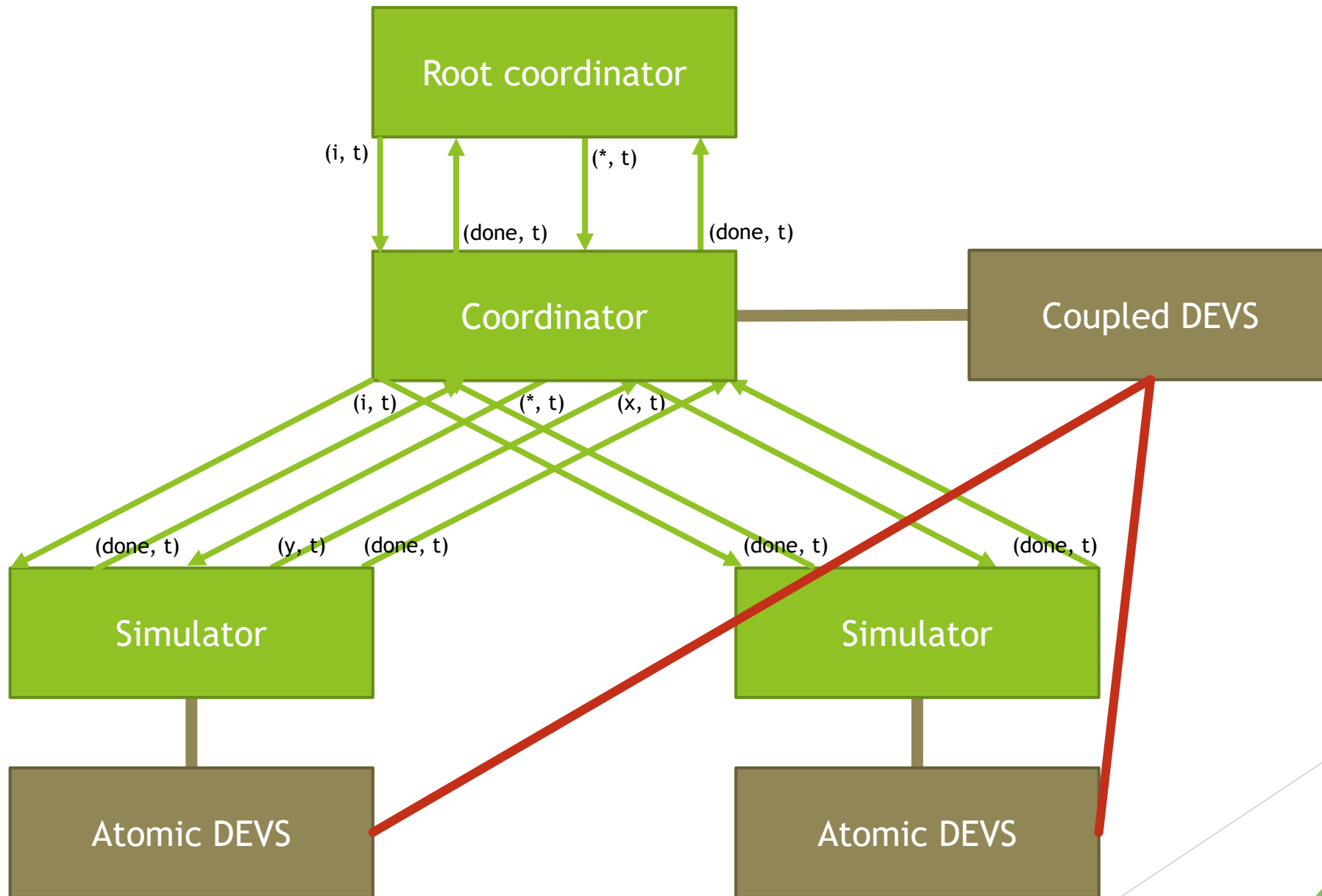
$$Z_{self,j} : X_{self} \rightarrow X_j, \forall j \in D$$

$$Z_{i,self} : Y_i \rightarrow Y_{self}, \forall i \in D$$

$$Z_{i,j} : Y_i \rightarrow X_j, \forall i, j \in D$$

$$select : 2^D \rightarrow D$$

$$\forall E \subseteq D, E \neq \emptyset : select(E) \in E$$



message m	simulator	coordinator
$(*, from, t)$	<p>simulator correct only if $t = t_N$</p> <p>$y \leftarrow \lambda(s)$</p> <p>if $y \neq \phi$:</p> <p> send $(\lambda(s), self, t)$ to parent</p> <p>$s \leftarrow \delta_{int}(s)$</p> <p>$t_L \leftarrow t$</p> <p>$t_N \leftarrow t_L + ta(s)$</p> <p>send $(done, self, t_N)$ to parent</p>	<p>send $(*, self, t)$ to i^*, where</p> <p>$i^* = select(imm_children)$</p> <p>$imm_children = \{i \in D \mid M_i.t_N = t\}$</p> <p>$active_children \leftarrow active_children \cup \{i^*\}$</p>

message m	simulator	coordinator
$(x, from, t)$	<p>simulator correct only if $t_L \leq t \leq t_N$ (ignore δ_{int} to resolve a $t = t_N$ conflict)</p> <p>$e \leftarrow t - t_L$</p> <p>$s \leftarrow \delta_{ext}(s, e, x)$</p> <p>$t_L \leftarrow t$</p> <p>$t_N \leftarrow t_L + ta(s)$</p> <p>send $(done, self, t_N)$ to parent</p>	<p>$\forall i \in I_{self} :$</p> <p>send $(Z_{self,i}(x), self, t)$ to i</p> <p>$active_children \leftarrow active_children \cup \{i\}$</p>

message m

simulator

coordinator

$(y, from, t)$

$\forall i \in I_{from} \setminus \{self\} :$

send $(Z_{from,i}(y), from, t)$ to i

$active_children \leftarrow active_children \cup \{i\}$

if $self \in I_{from} :$

send $(Z_{from,self}(y), self, t)$ to parent

$(done, from, t)$

$active_children \leftarrow active_children \setminus \{from\}$

if $active_children = \emptyset :$

$t_L \leftarrow t$

$t_N \leftarrow \min\{M_i.t_N \mid i \in D\}$

send $(done, self, t_N)$ to parent

$t \leftarrow t_N$ of topmost coordinator

repeat until $t \geq t_{end}$ (or some other termination condition)

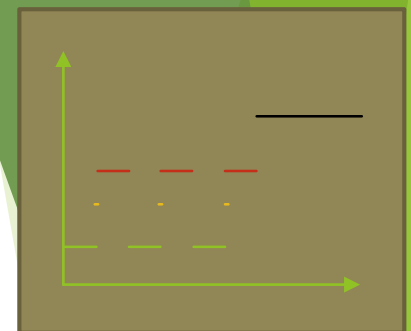
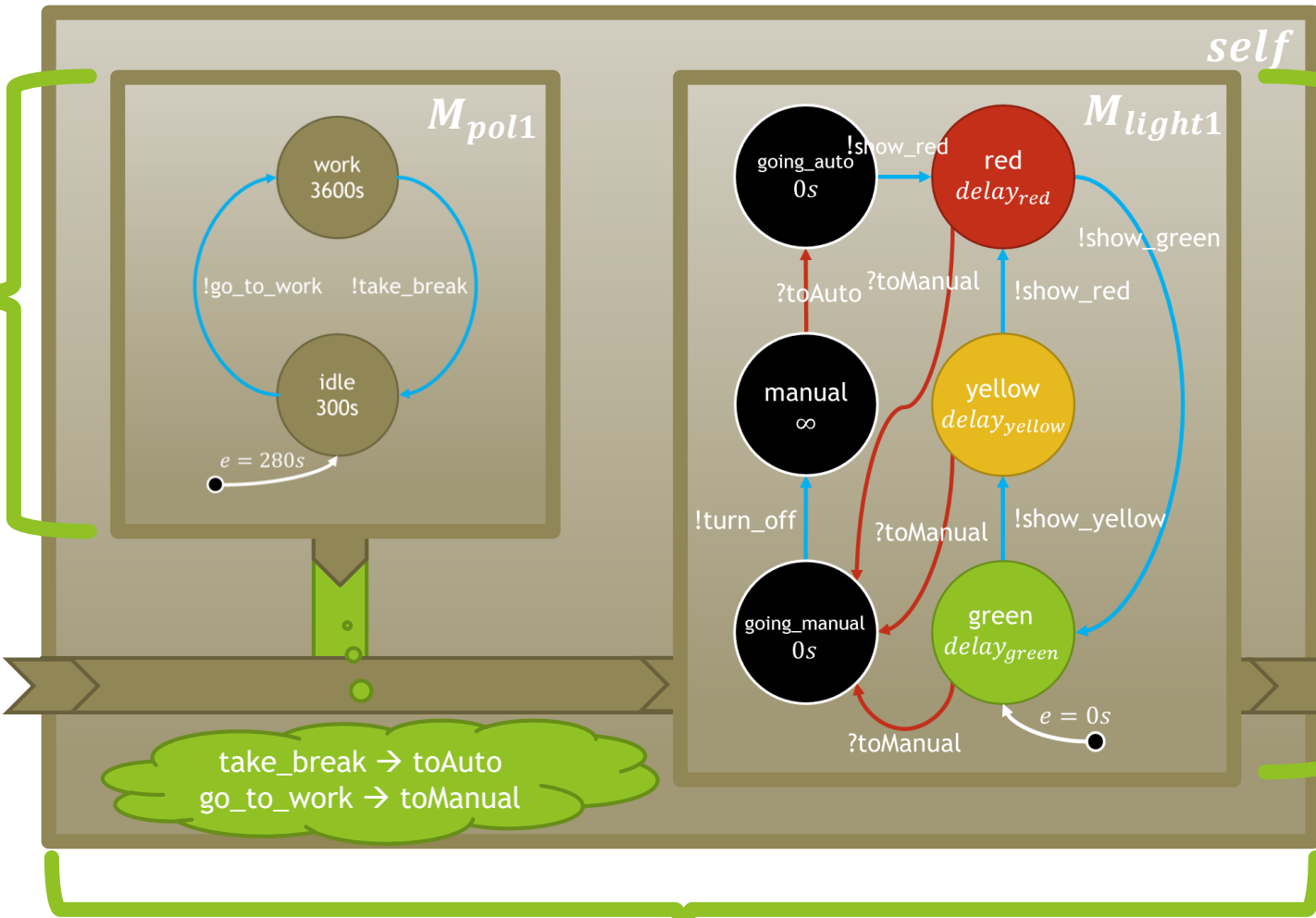
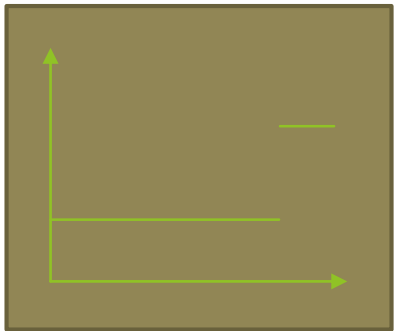
send $(*, main, t)$ to topmost coupled model top

wait for $(done, top, t_N)$

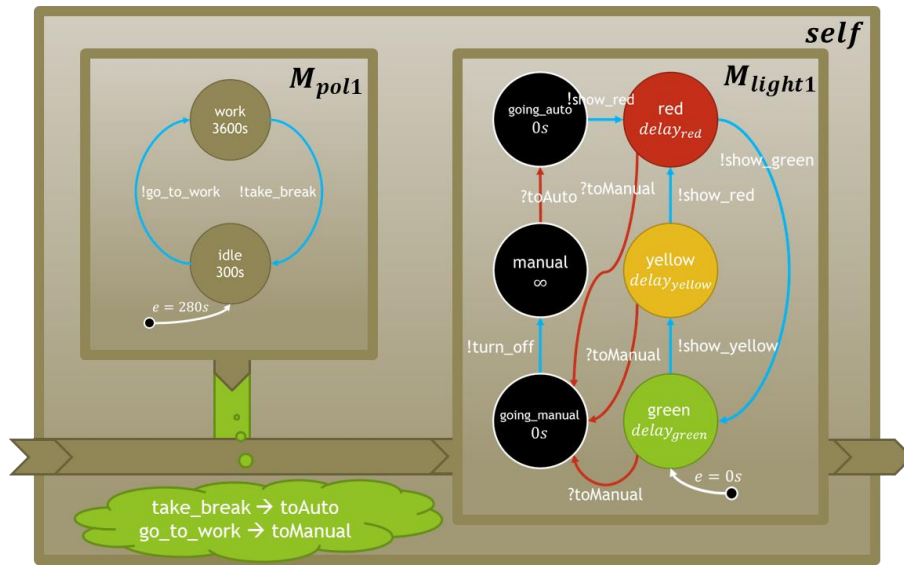
$t \leftarrow t_N$

Closure Under Coupling

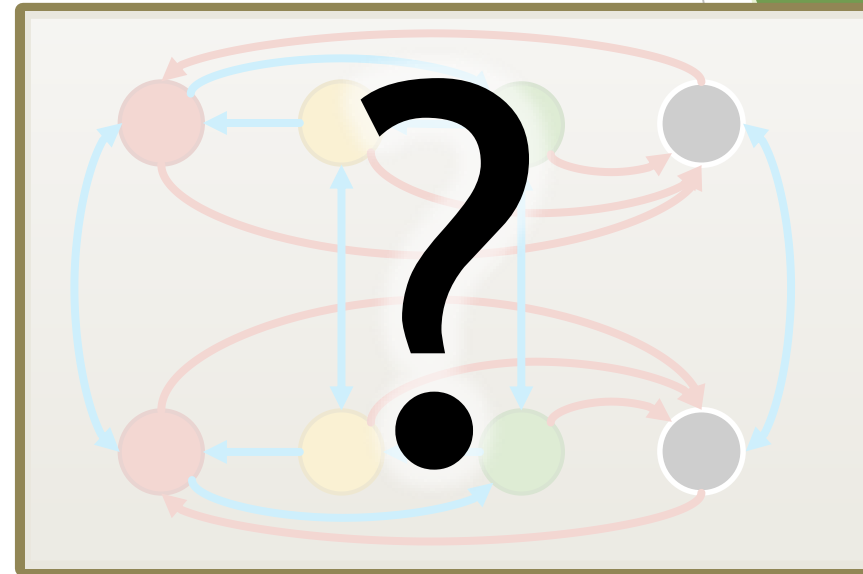
Denotational semantics for Coupled DEVS models



?

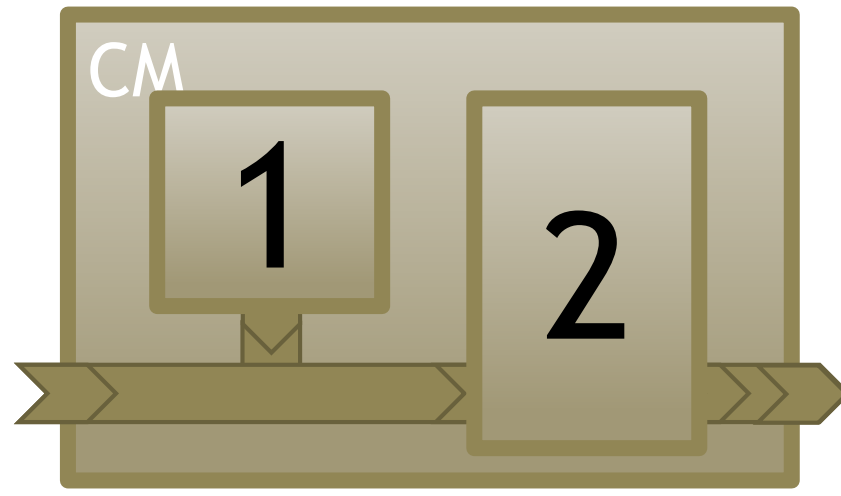


flatten

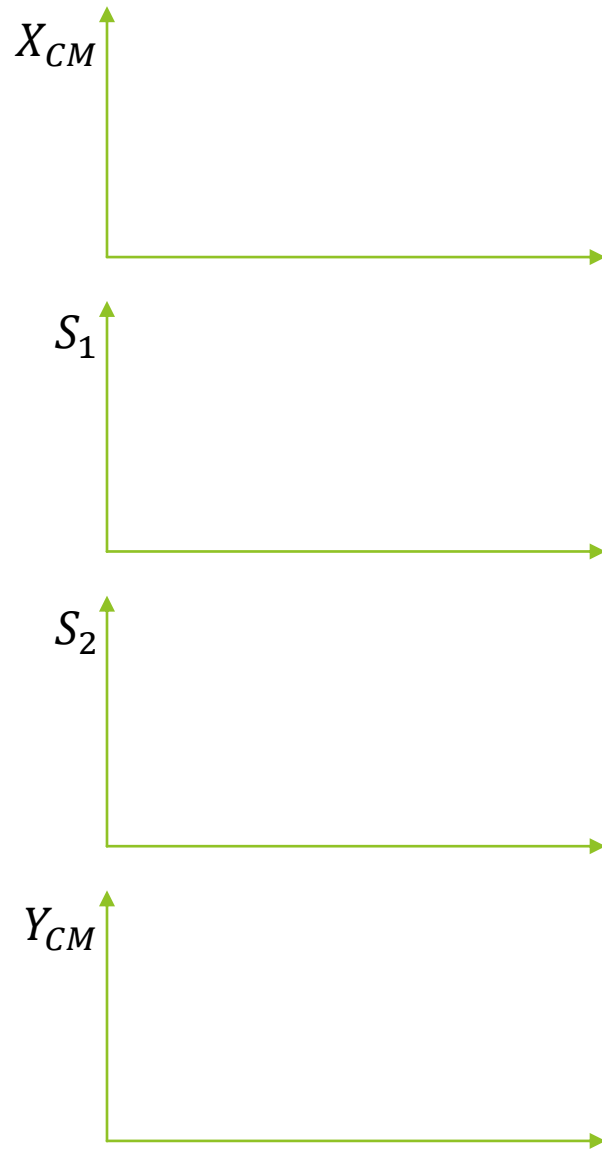


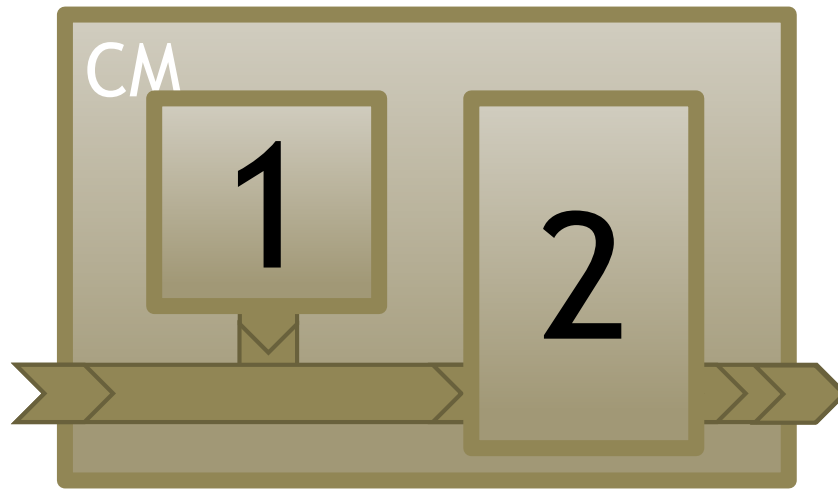
$$CM = \langle X_{self}, Y_{self}, D, MS, IS, ZS \rangle$$

$$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$



$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

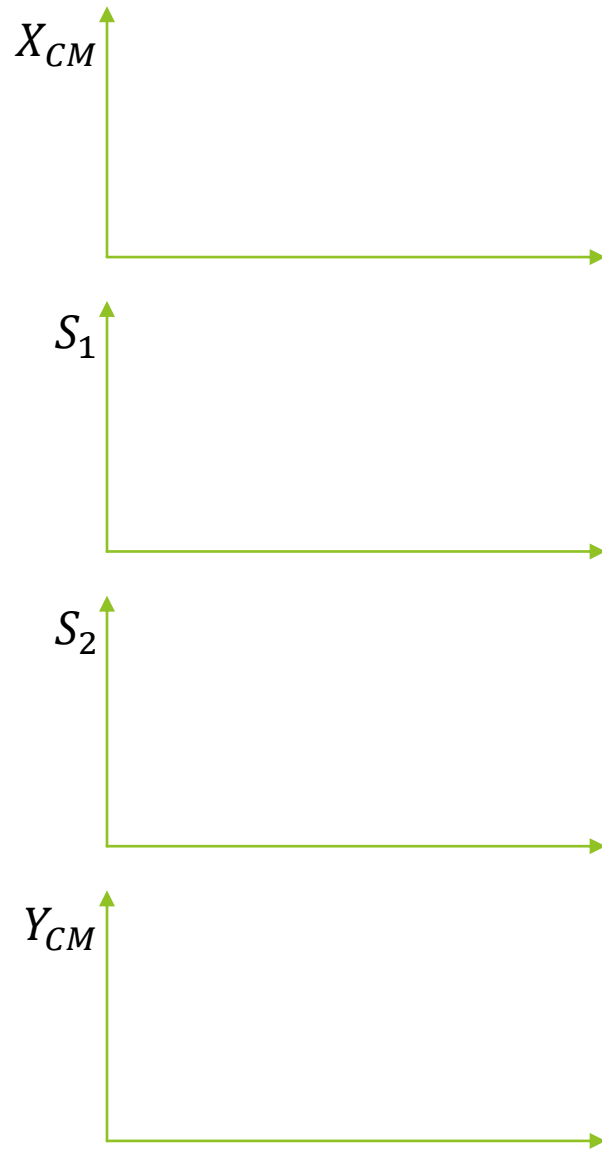


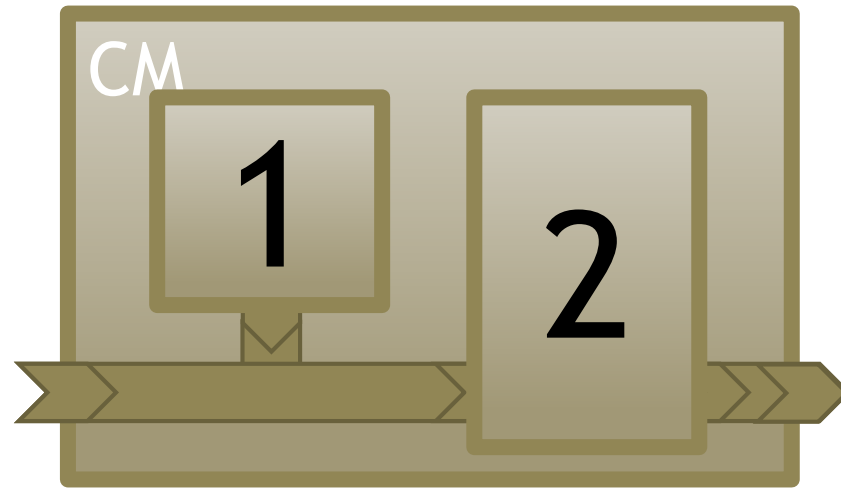
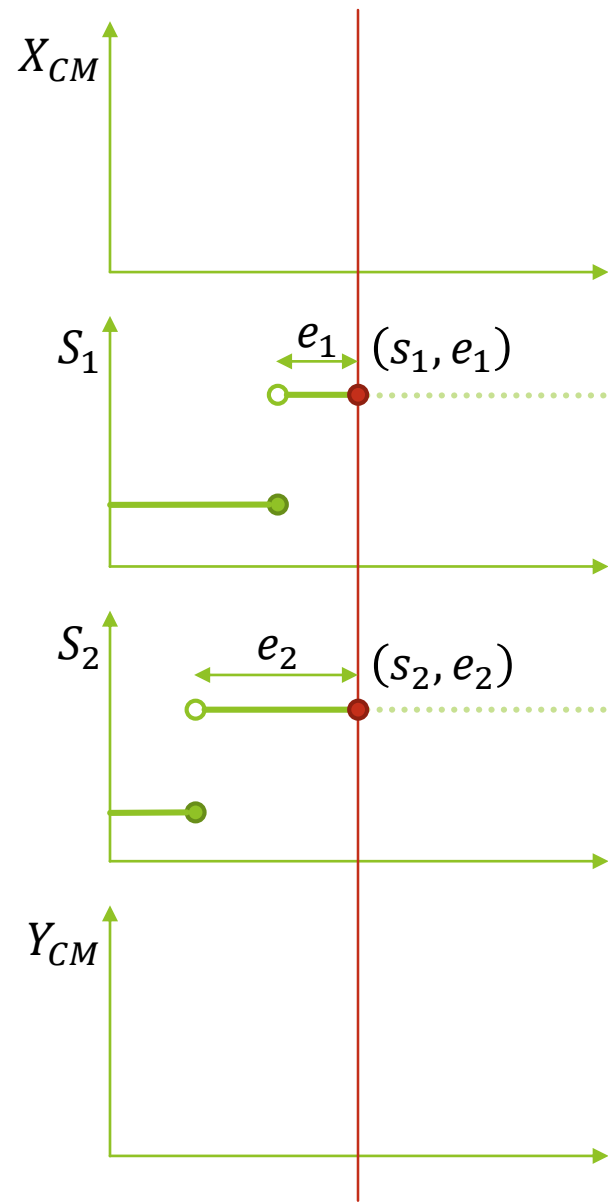


$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

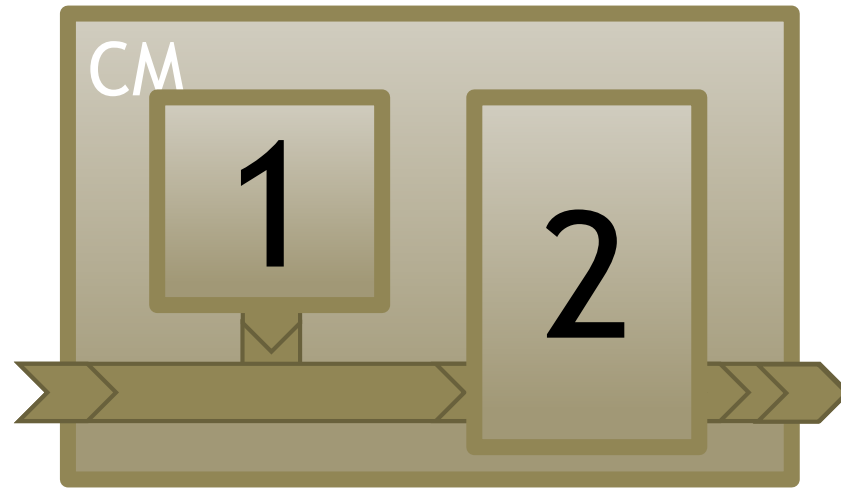
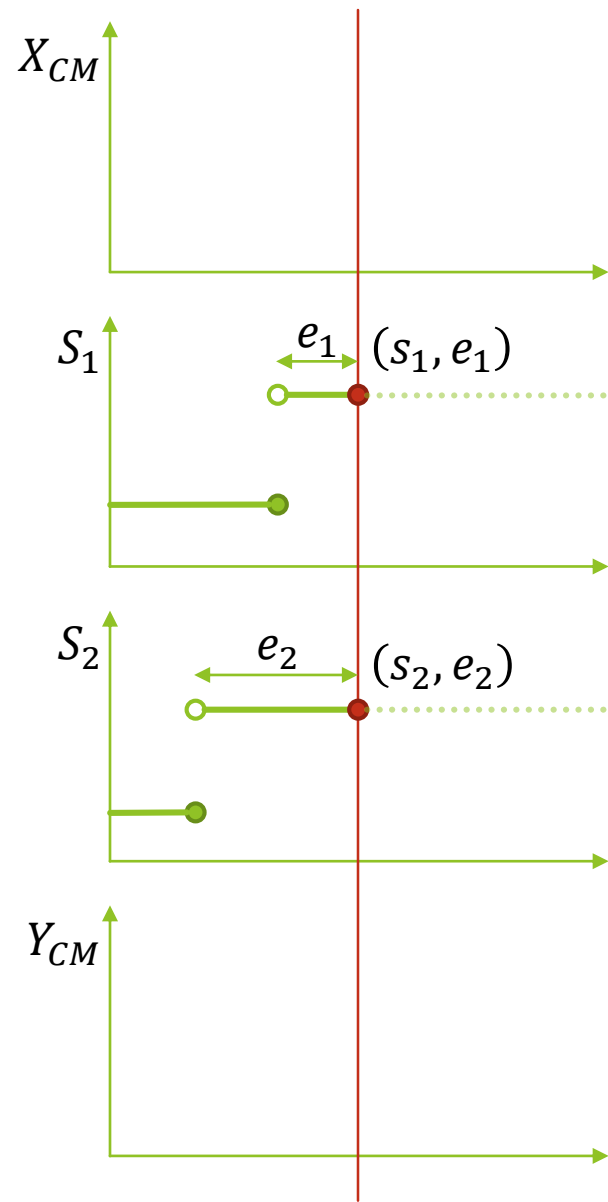
$$X = X_{CM}$$

$$Y = Y_{CM}$$



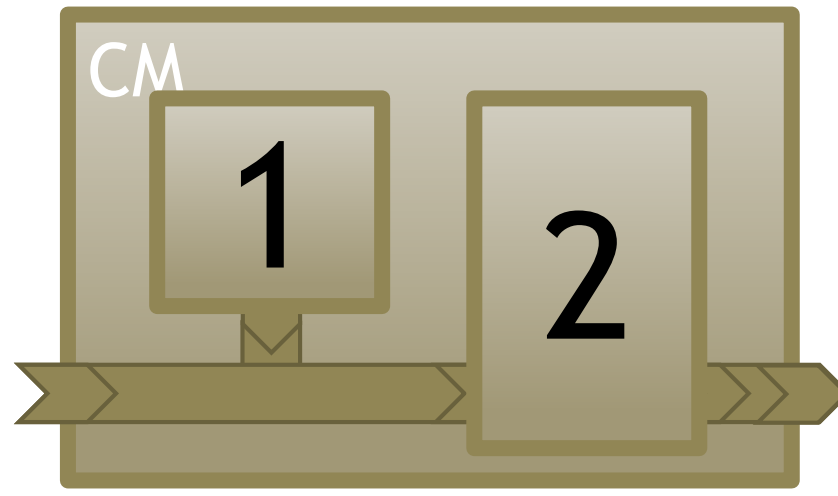


$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

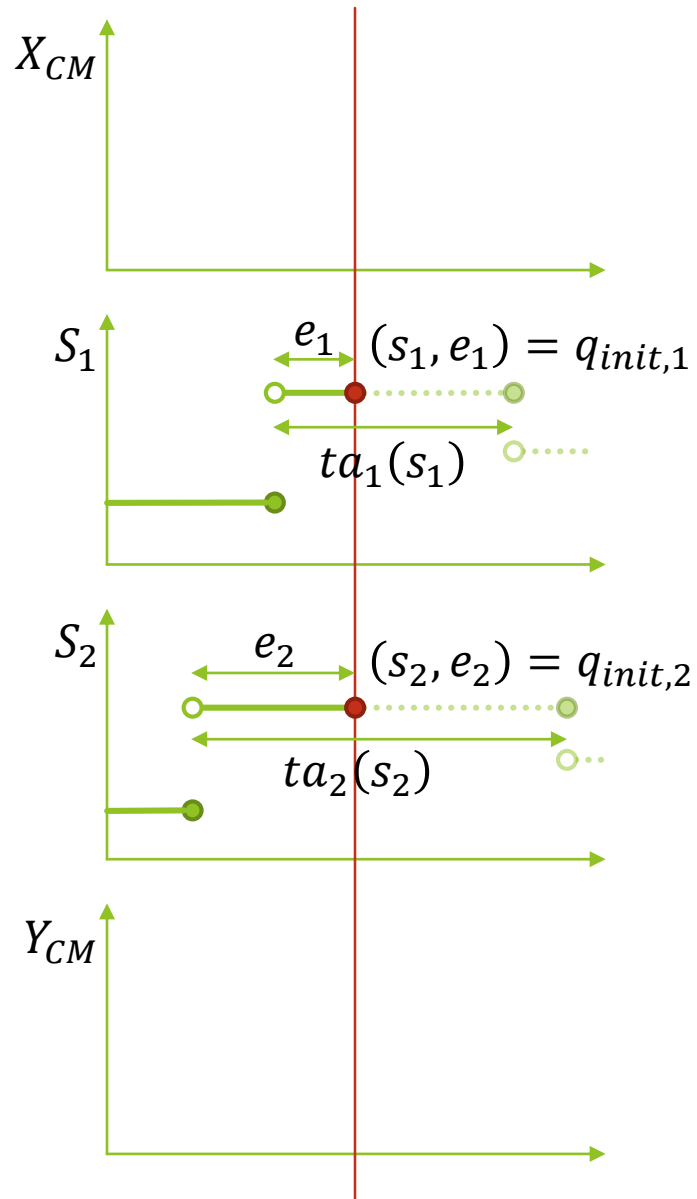


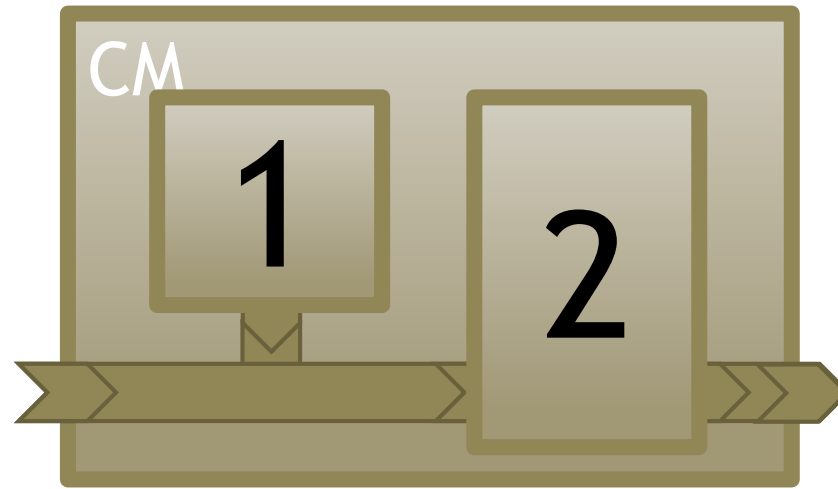
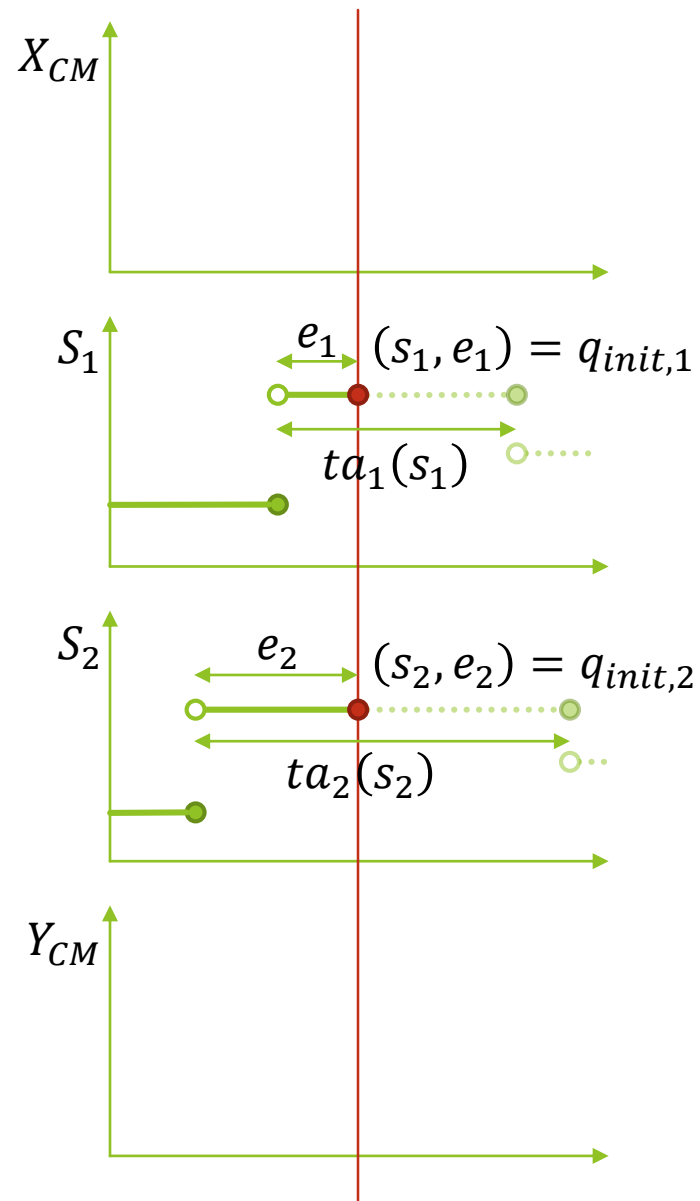
$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

$$S = \times_{i \in D} Q_i$$



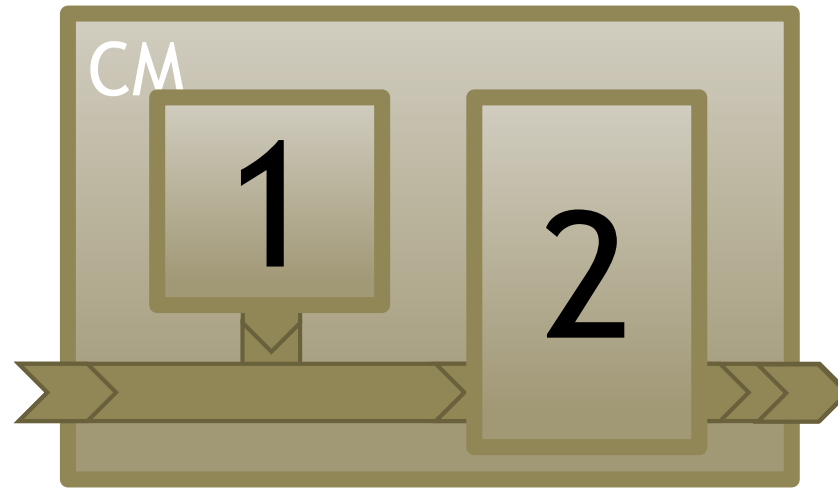
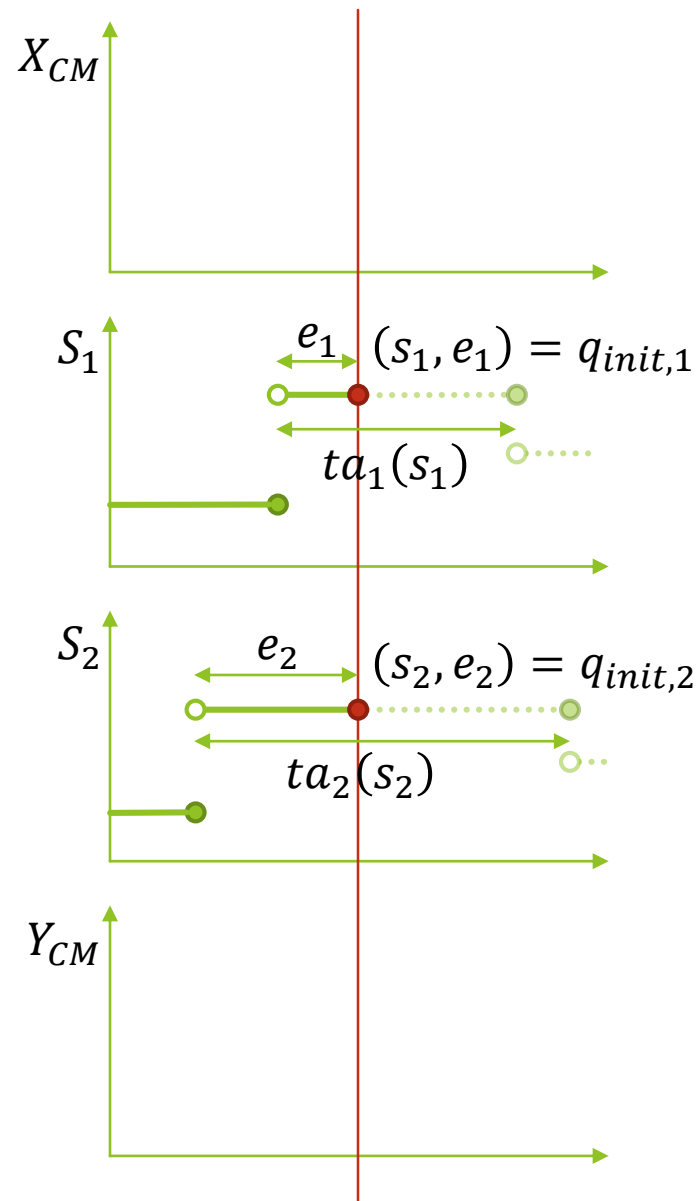
$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$





$$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

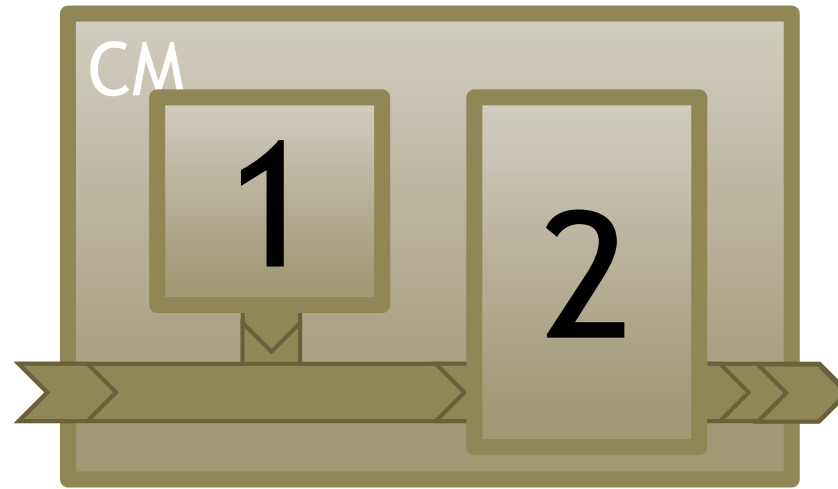
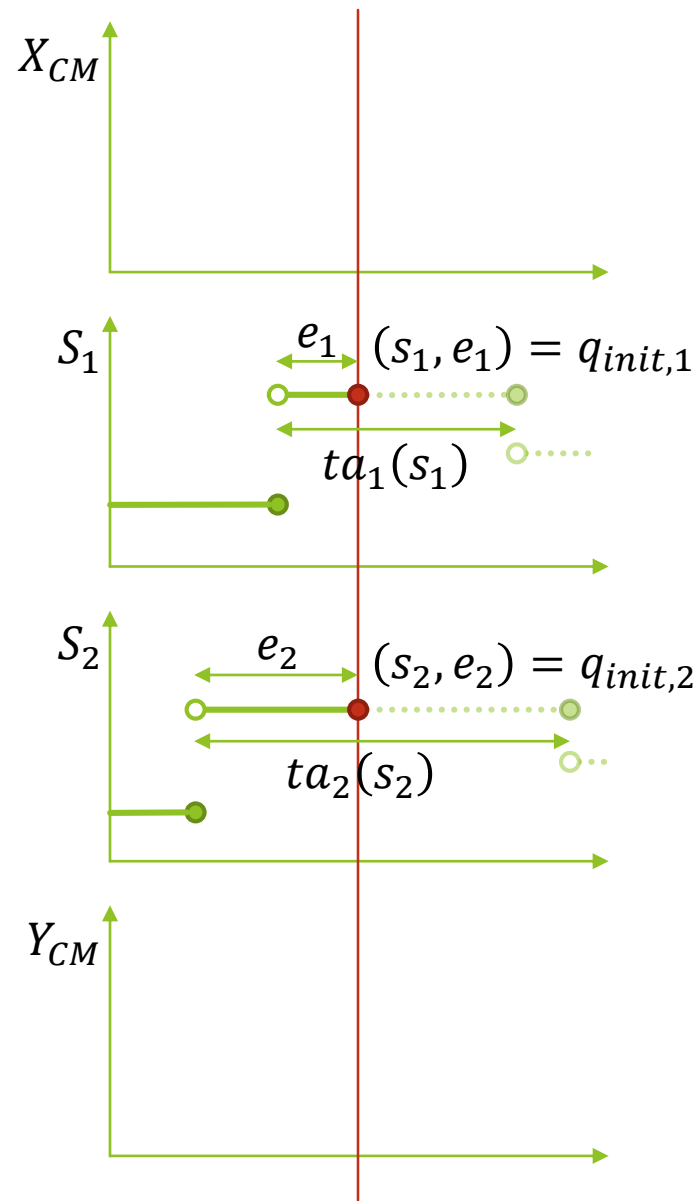
$$q_{init} = (s_{init}, e_{init})$$



$$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

$$q_{init} = (s_{init}, e_{init})$$

$$s_{init} = (\dots, (s_{init,i}, e_{init,i} - e_{init}), \dots)$$

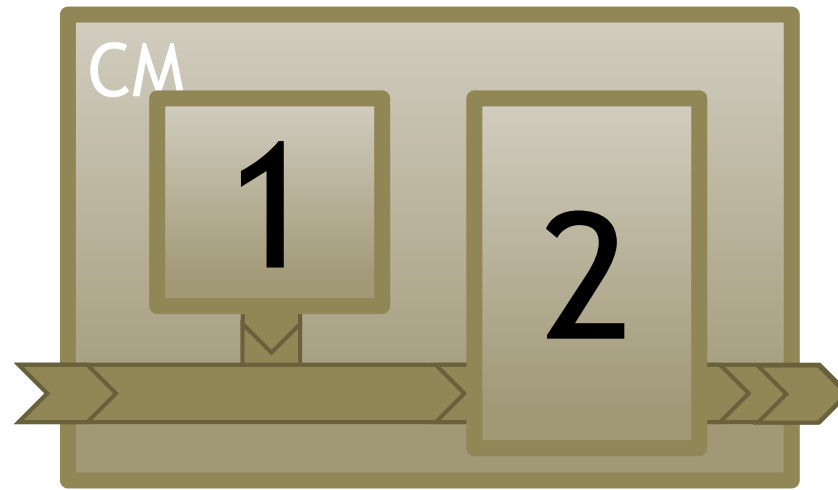


$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

$$q_{init} = (s_{init}, e_{init})$$

$$s_{init} = (\dots, (s_{init,i}, e_{init,i} - e_{init}), \dots)$$

$$e_{init} = \min_{i \in D} \{e_{init,i}\}$$



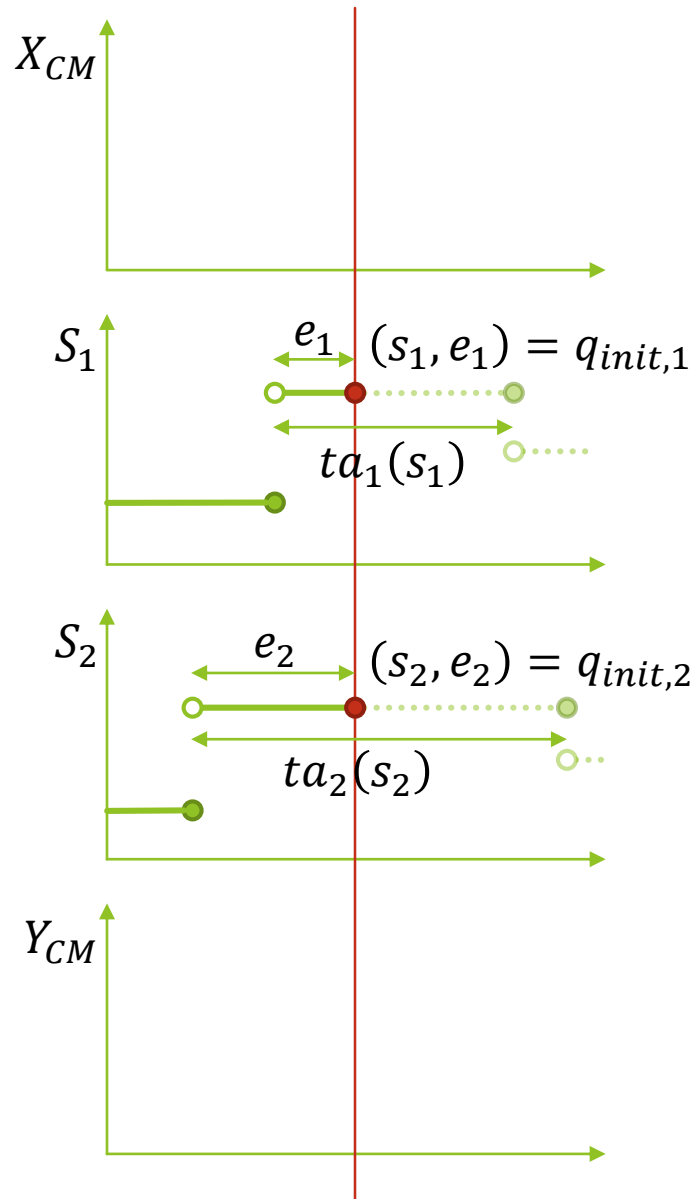
$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

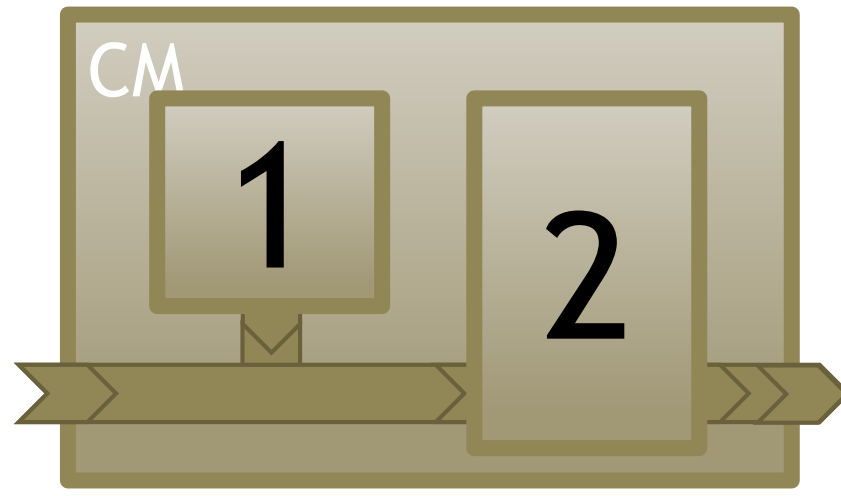
$$q_{init} = (s_{init}, e_{init})$$

$$s_{init} = (\dots, (s_{init,i}, e_{init,i} - e_{init}), \dots)$$

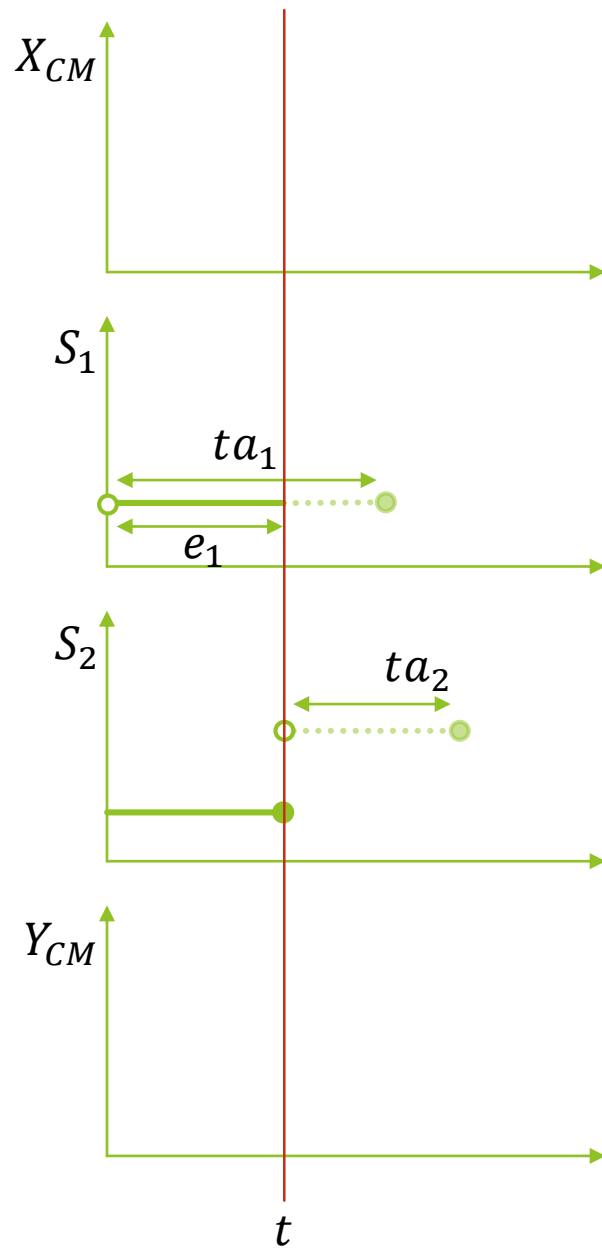
$$e_{init} = \min_{i \in D} \{e_{init,i}\}$$

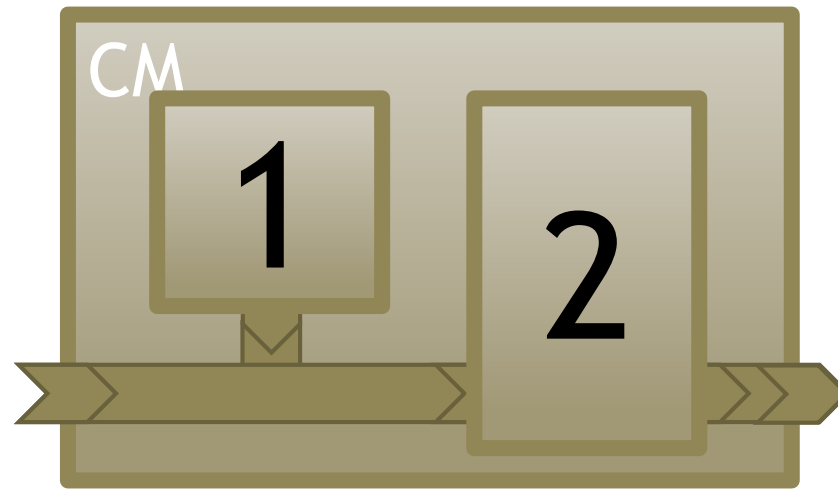
$$(s_{init,i}, e_{init,i}) = q_{init,i}$$





$$\text{flatten}(CM) = \langle X, Y, S, q_{\text{init}}, \delta_{\text{int}}, \delta_{\text{ext}}, \lambda, ta \rangle$$

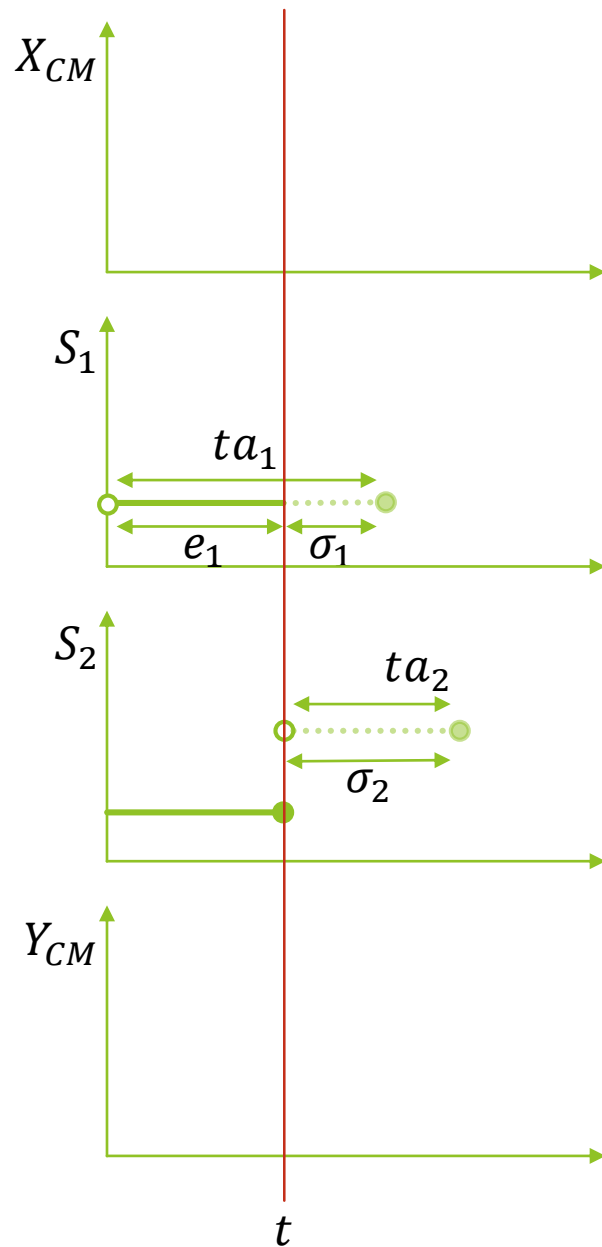


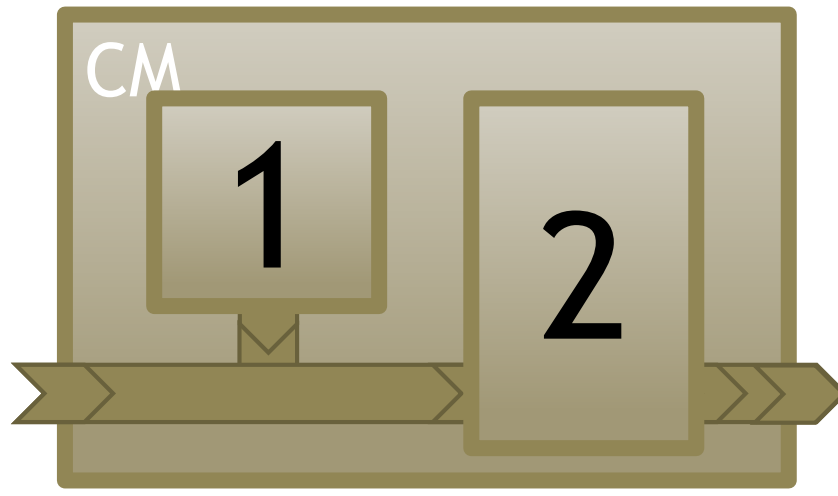


$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

$$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$$

$$ta(s) = \min_{i \in D} \{ \sigma_i = ta_i(s_i) - e_i \}$$





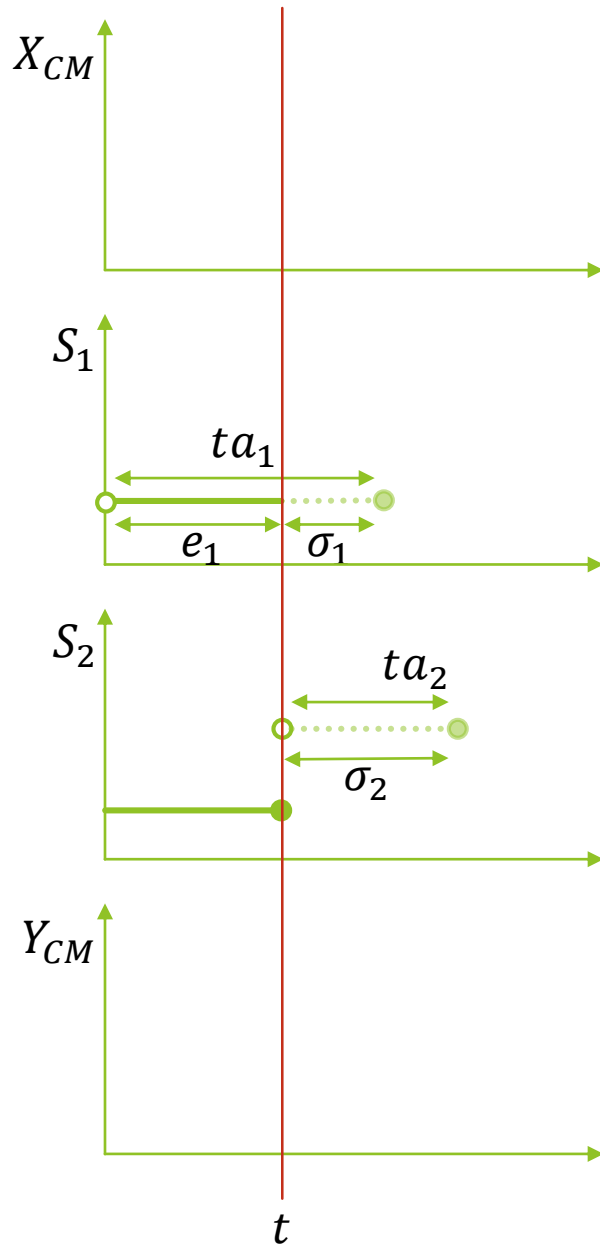
$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

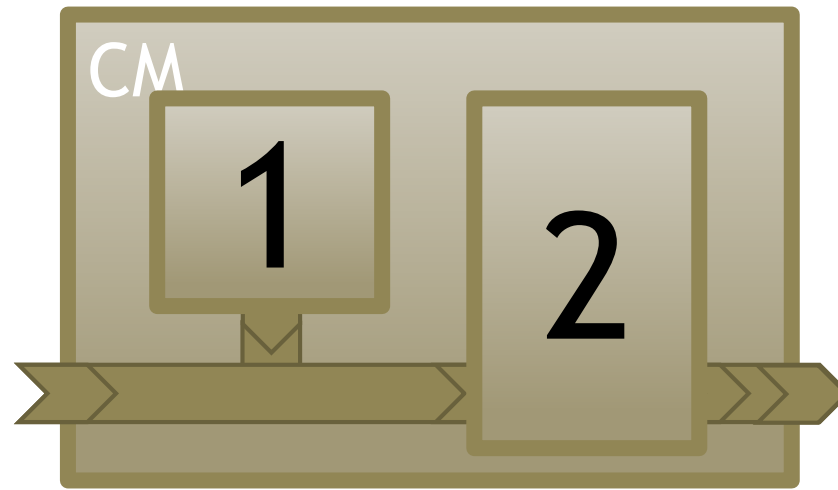
$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$

$ta(s) = \min_{i \in D} \{ \sigma_i = ta_i(s_i) - e_i \}$

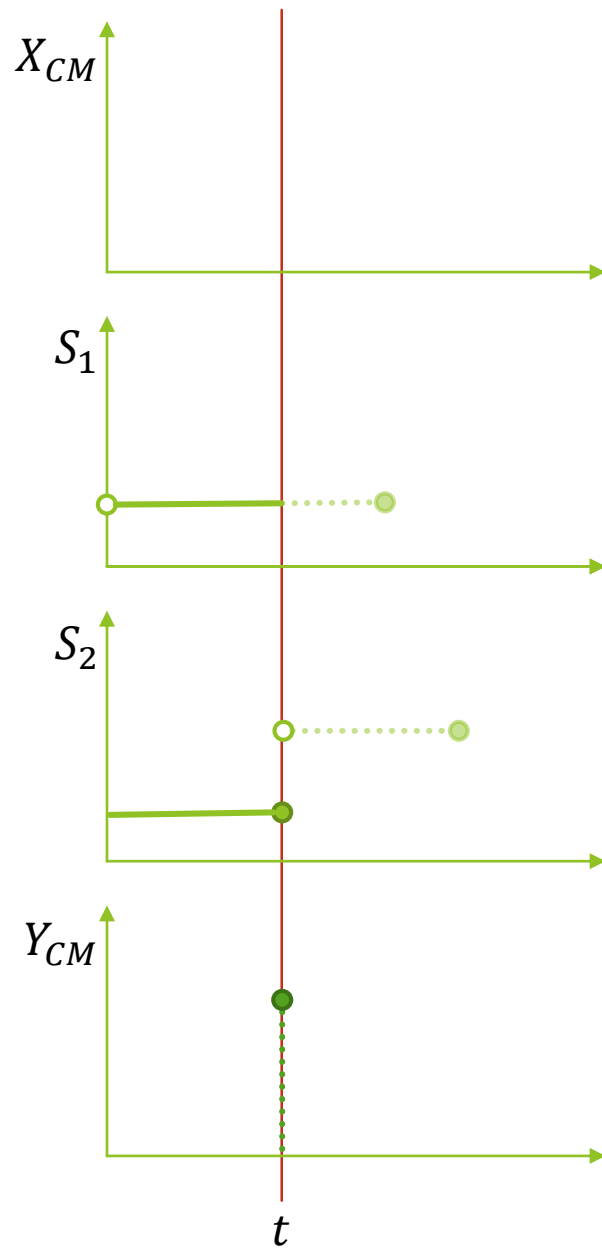
$IMM(s) = \{ i \in D \mid \sigma_i = ta(s) \}$

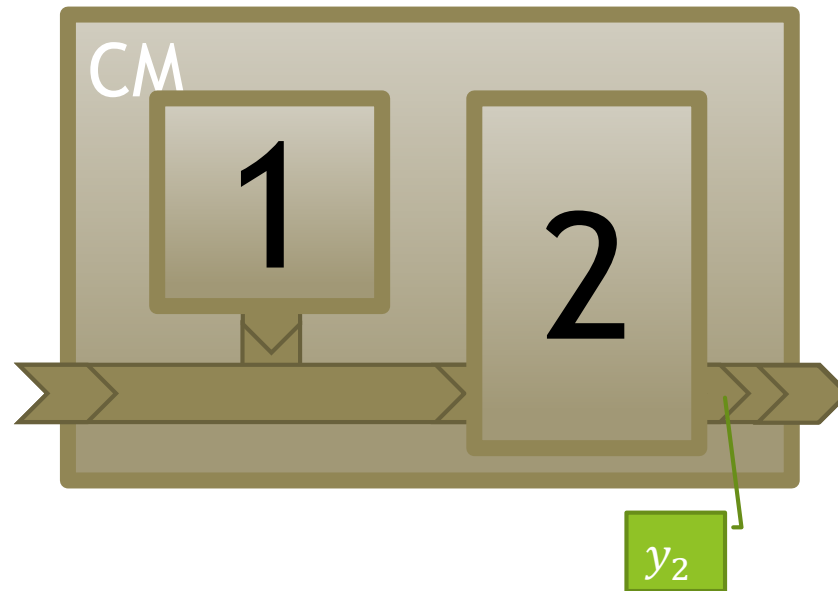
$select(IMM(s)) = i^*$





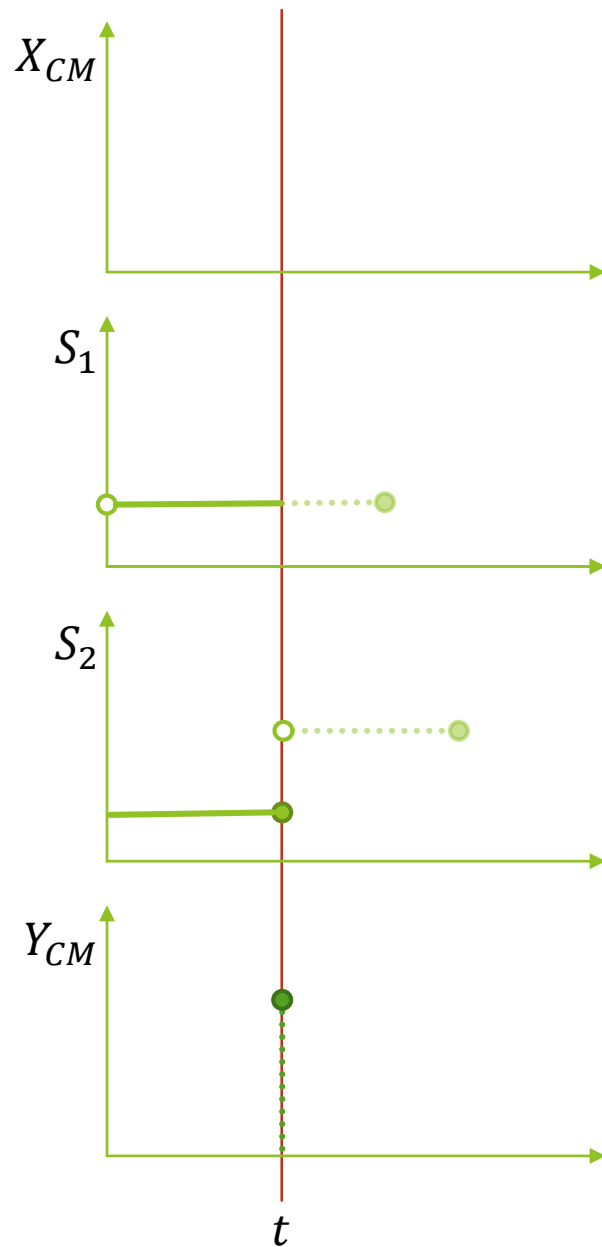
$$\text{flatten}(CM) = \langle X, Y, S, q_{\text{init}}, \delta_{\text{int}}, \delta_{\text{ext}}, \lambda, ta \rangle$$

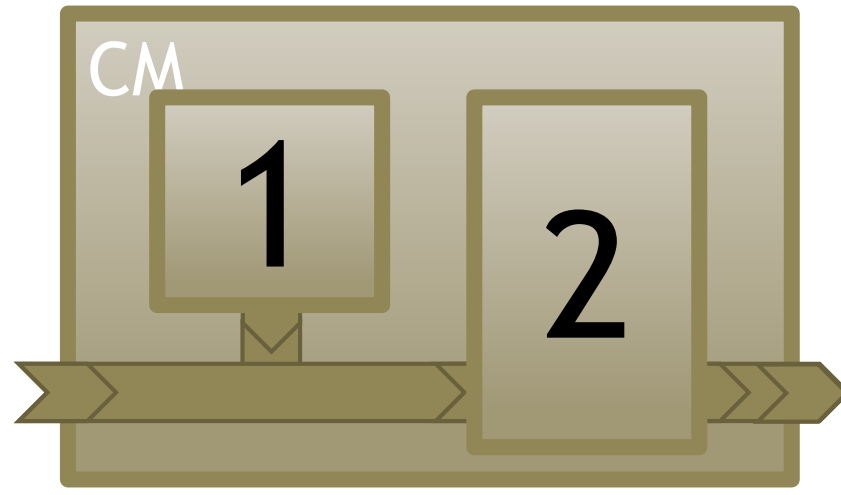
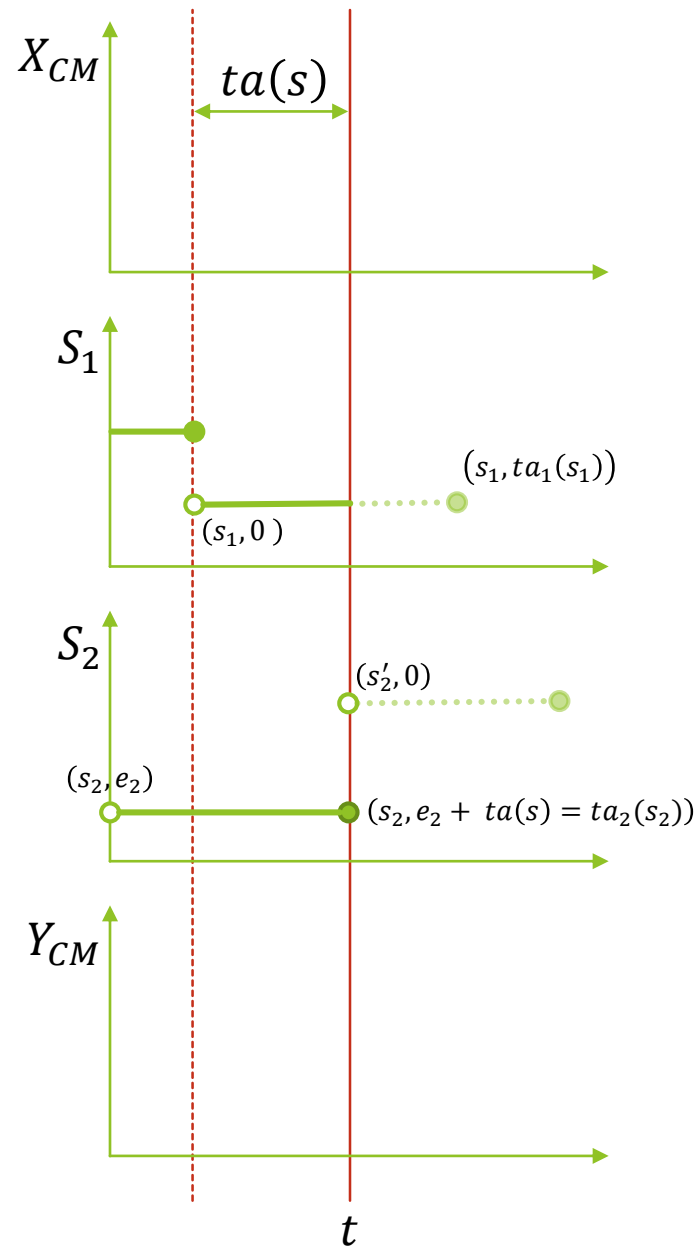




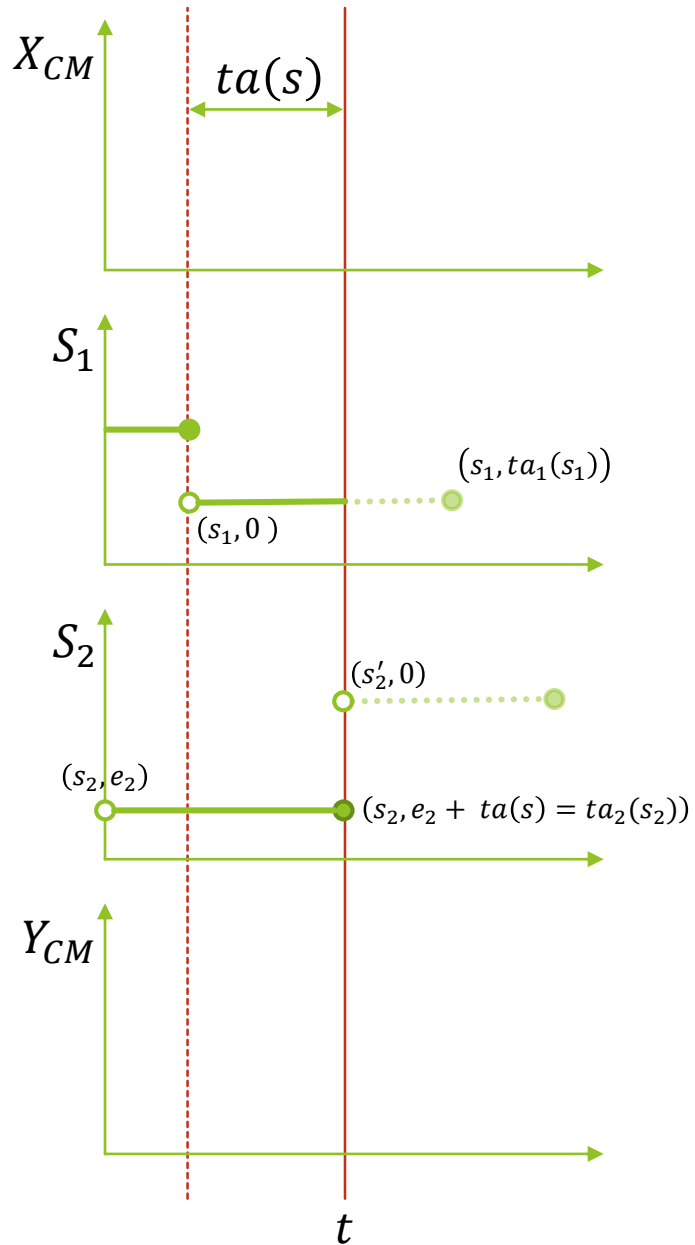
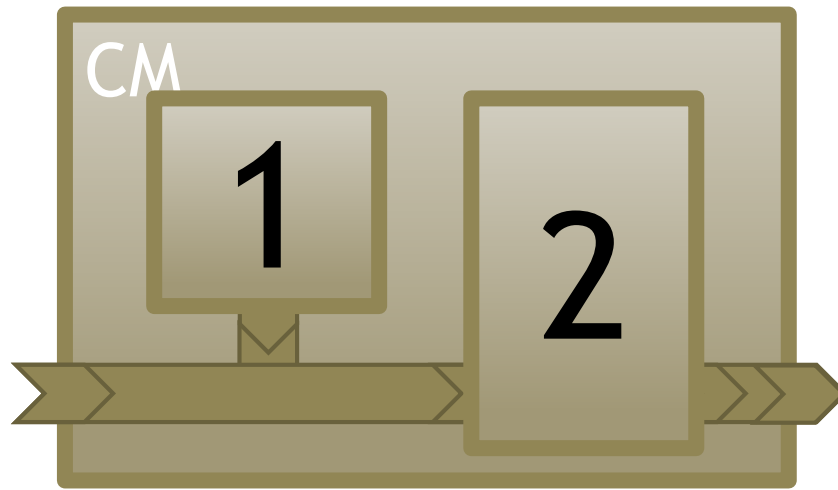
$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

$$\lambda(s) = \begin{cases} Z_{i^*, self}(\lambda_{i^*}(s_{i^*})) & \text{if } self \in I_{i^*} \\ \emptyset & \text{if } self \notin I_{i^*} \end{cases}$$





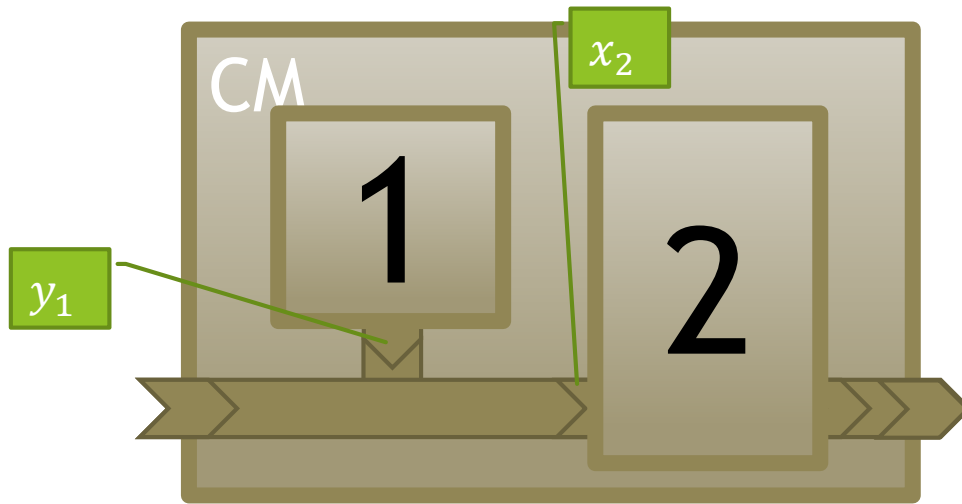
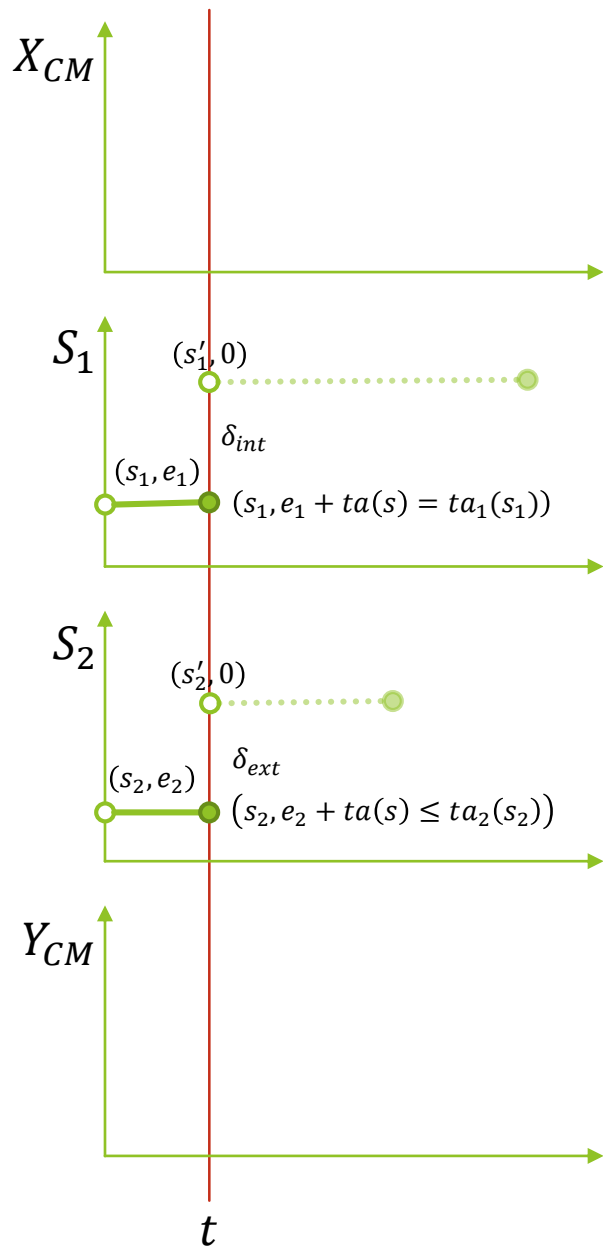
$$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$



$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

$$\delta_{int}(s) = (\dots, (s'_j, e'_j), \dots)$$

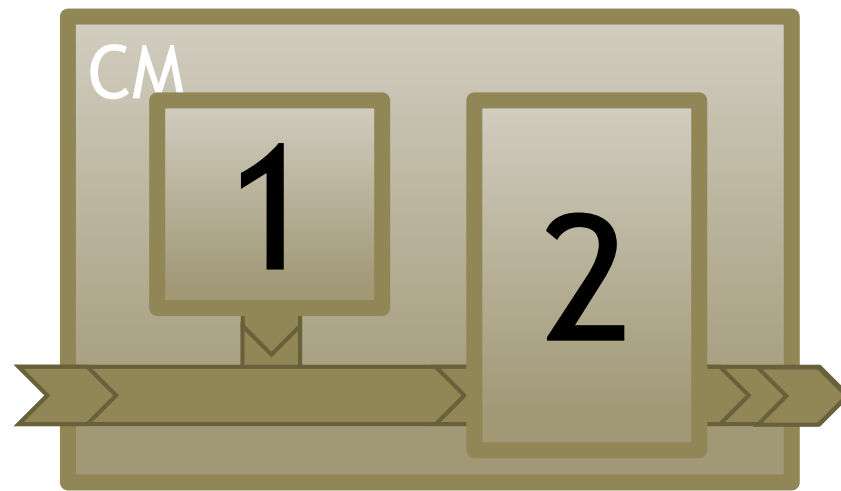
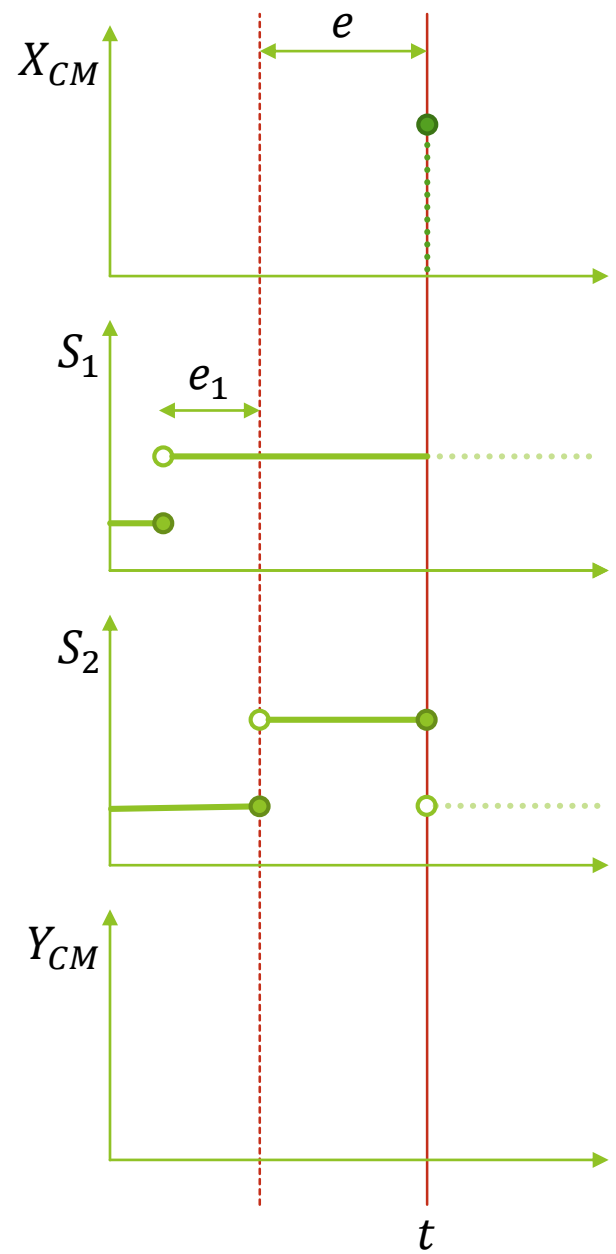
$$(s'_j, e'_j) = \begin{cases} (\delta_{int,j}(s_j), 0) & \text{for } j = i^* \\ ? & \text{for } j \in I_{i^*} \setminus \{self\} \\ (s_j, e_j + ta(s)) & \text{else} \end{cases}$$



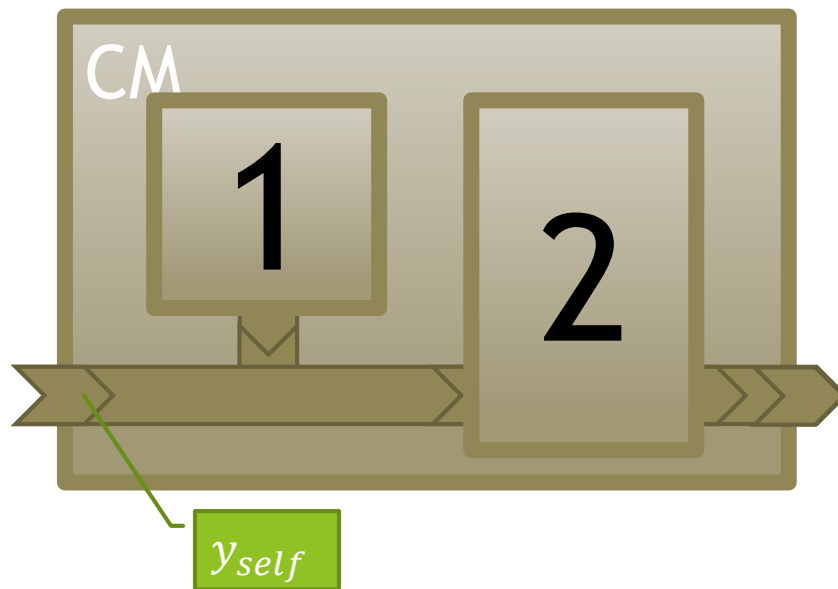
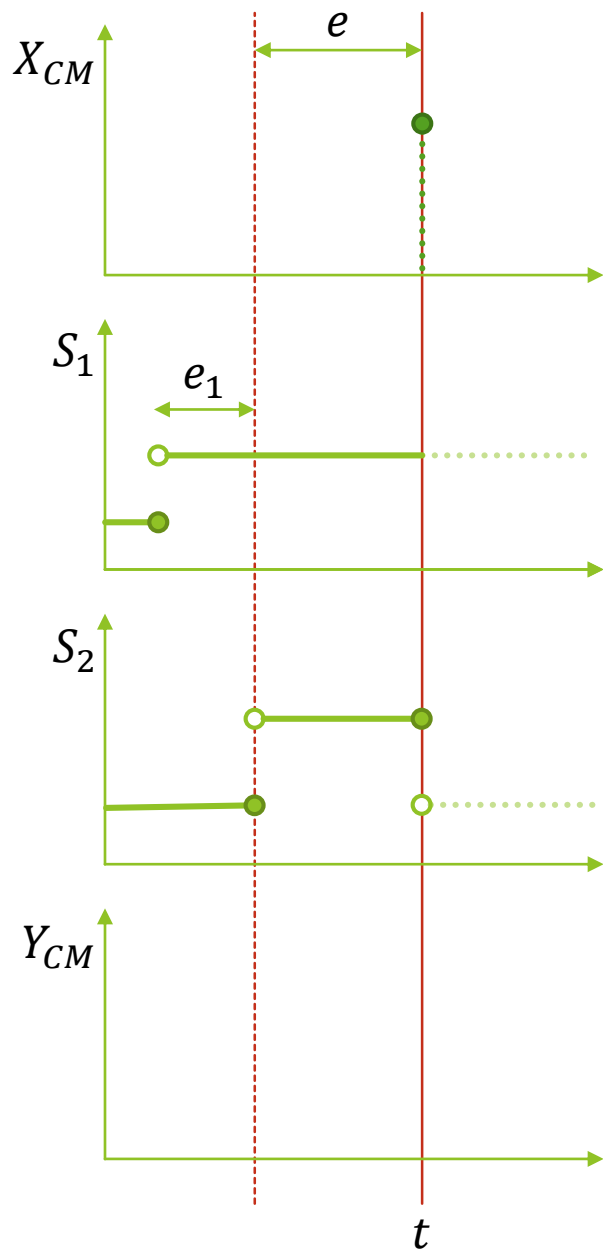
$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

$$\delta_{int}(s) = (\dots, (s'_j, e'_j), \dots)$$

$$(s'_j, e'_j) = \begin{cases} (\delta_{int,j}(s_j), 0) & \text{for } j = i^* \\ \left(\delta_{ext,j} \left((s_j, e_j + ta(s)), Z_{i^*,j}(\lambda_{i^*}(s_{i^*})) \right), 0 \right) & \text{for } j \in I_{i^*} \setminus \{\text{self}\} \\ (s_j, e_j + ta(s)) & \text{else} \end{cases}$$



$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$



$flatten(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$

$\delta_{ext}((s, e), x) = (\dots, (s'_i, e'_i), \dots)$

$$(s'_i, e'_i) = \begin{cases} (\delta_{ext,i}((s_i, e_i + e), Z_{self,i}(x)), 0) & \text{for } i \in I_{self} \\ (s_i, e_i + e) & \text{else} \end{cases}$$

$$\text{flatten}(CM) = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

$$X = X_{CM}$$

$$Y = Y_{CM}$$

$$S = \times_{i \in D} Q_i$$

$$Q = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$$

$$q_{init} = (s_{init}, e_{init}) \in Q$$

$$s_{init} = (\dots, (s_{init,i}, e_{init,i} - e_{init}), \dots)$$

$$e_{init} = \min_{i \in D} \{e_{init,i}\}$$

$$(s_{init,i}, e_{init,i}) = q_{init,i}$$

$$\delta_{int}(s) = (\dots, (s'_j, e'_j), \dots)$$

$$(s'_j, e'_j) = \begin{cases} (\delta_{int,j}(s_j), 0) & \text{for } j = i^* \\ \left(\delta_{ext,j} \left((s_j, e_j + ta(s)), Z_{i^*,j}(\lambda_{i^*}(s_{i^*})) \right), 0 \right) & \text{for } j \in I_{i^*} \\ (s_j, e_j + ta(s)) & \text{else} \end{cases}$$

$$\delta_{ext}((s, e), x) = (\dots, (s'_i, e'_i), \dots)$$

$$(s'_i, e'_i) = \begin{cases} \left(\delta_{ext,i} \left((s_i, e_i + e), Z_{self,i}(x) \right), 0 \right) & \text{for } i \in I_{self} \\ (s_i, e_i + e) & \text{else} \end{cases}$$

$$\lambda(s) = \begin{cases} Z_{i^*,self}(\lambda_{i^*}(s_{i^*})) & \text{if } self \in I_{i^*} \\ \phi & \text{if } self \notin I_{i^*} \end{cases}$$

$$i^* = \text{select}(IMM(s))$$

$$IMM(s) = \{i \in D \mid \sigma_i = ta(s)\}$$

$$ta(s) = \min_{i \in D} \{\sigma_i = ta_i(s_i) - e_i\}$$

Applications of DEVS

DEVS in practice

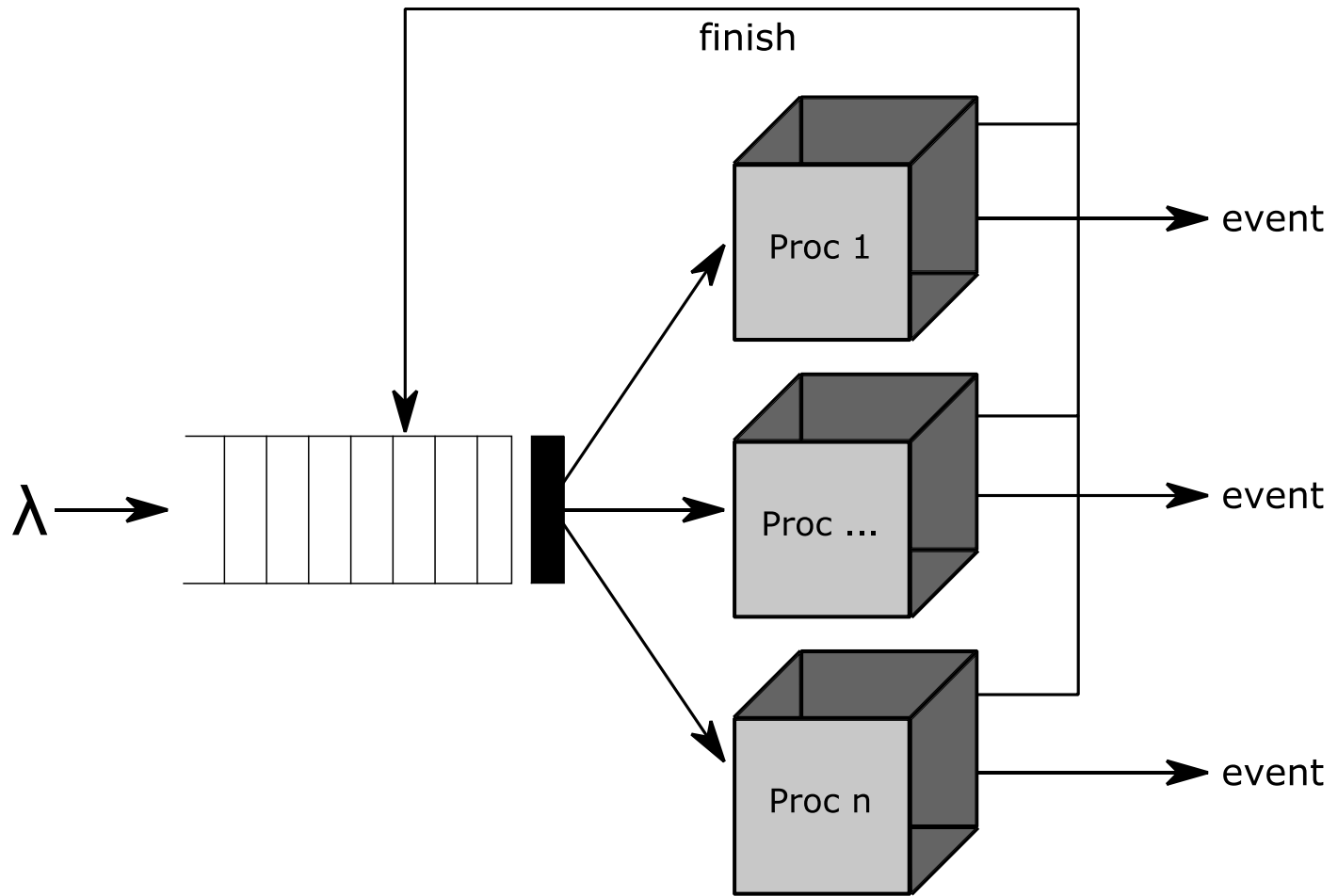


Table 3. M/M/1 Queueing System

$$\rho = \lambda W_s, \quad p_n = P[N = n] = (1 - \rho)\rho^n, \quad n = 0, 1, \dots$$

$$P[N \geq n] = \rho^n, \quad n = 0, 1, \dots$$

$$L = E[N] = \lambda W = \frac{\rho}{1 - \rho}, \quad \sigma_N^2 = \frac{\rho}{(1 - \rho)^2}.$$

$$L_q = \lambda W_q = \frac{\rho^2}{1 - \rho}, \quad \sigma_{N_q}^2 = \frac{\rho^2(1 + \rho - \rho^2)}{(1 - \rho)^2}.$$

$$E[N_q | N_q > 0] = \frac{1}{1 - \rho}, \quad \text{Var}[N_q | N_q > 0] = \frac{\rho}{(1 - \rho)^2}.$$

$$W[t] = P[w \leq t] = 1 - \exp\left(\frac{-t}{W}\right), \quad P[w > t] = \exp\left(\frac{-t}{W}\right).$$

$$W = E[w] = \frac{W_s}{1 - \rho}, \quad \sigma_w^2 = W^2.$$

$$\pi_w[r] = W \ln\left(\frac{100}{100 - r}\right), \quad \pi_w[90] = W \ln 10, \quad \pi_w[95] = W \ln 20$$

$$W_q[t] = P[q \leq t] = 1 - \rho \exp\left(\frac{-t}{W}\right), \quad P[q > t] = \rho \exp\left(\frac{-t}{W}\right).$$

$$W_q = \frac{\rho W_s}{1 - \rho}, \quad \sigma_q^2 = \frac{(2 - \rho)\rho W_s^2}{(1 - \rho)^2}.$$

$$\pi_q[r] = \max\left\{W \ln\left(\frac{100\rho}{100 - r}\right), 0\right\}.$$

$$\pi_q[90] = \max\{W \ln(10\rho), 0\}, \quad \pi_q[95] = \max\{W \ln(20\rho), 0\}.$$

Table 4. M/M/1/K Queueing System

$$p_n = \begin{cases} \frac{(1 - a)a^n}{(1 - a^{K+1})} & \text{if } \lambda \neq \mu, \\ \frac{1}{K + 1} & \text{if } \lambda = \mu, \end{cases}$$

$n = 0, 1, \dots, K$, where $a = \lambda W_s$.

$\lambda_a = (1 - p_K)\lambda$, Mean arrival rate into system.

$$L = \begin{cases} \frac{a[1 - (K + 1)a^K + Ka^{K+1}]}{(1 - a)(1 - a^{K+1})} & \text{if } \lambda \neq \mu, \\ \frac{K}{2} & \text{if } \lambda = \mu. \end{cases}$$

$$L_q = L - (1 - p_0), \quad q_n = \frac{p_n}{1 - p_K}, \quad n = 0, 1, \dots, K - 1.$$

$$W[t] = 1 - \sum_{n=0}^{K-1} q_n Q[n; \mu t],$$

where

$$q[n; \mu t] = e^{-\mu t} \sum_{k=0}^n \frac{\mu t^k}{k!}.$$

$$W = \frac{L}{\lambda_a}, \quad W_q = \frac{L_q}{\lambda_a}.$$

$$W_q[t] = 1 - \sum_{n=0}^{K-2} q_{n+1} Q[n; \mu t].$$

$$E[q | q > 0] = \frac{W_q}{1 - p_0}, \quad \rho = (1 - p_K)a.$$

Table 3. M/M/1 Queueing System

$$\rho = \lambda W_s, \quad p_n = P[N = n] = (1 - \rho)\rho^n, \quad n = 0, 1, \dots$$

$$P[N \geq n] = \rho^n, \quad n = 0, 1, \dots$$

$$L = E[N] = \lambda W = \frac{\rho}{1 - \rho}, \quad \sigma_N^2 = \frac{\rho}{(1 - \rho)^2}$$

$$L_q = \lambda W_q = \frac{\rho^2}{1 - \rho}, \quad \sigma_{N_q}^2 = \frac{\rho^2(1 + \rho - \rho^2)}{(1 - \rho)^2}$$

$$E[N_q | N_q > 0] = \frac{1}{1 - \rho}, \quad \text{Var}[N_q | N_q > 0] = \frac{\rho}{(1 - \rho)^2}$$

$$W[t] = P[w \leq t] = 1 - \exp\left(\frac{-t}{W}\right), \quad P[w > t] = \exp\left(\frac{-t}{W}\right)$$

$$W = E[w] = \frac{W_s}{1 - \rho}$$

$$\pi_w[r] = \frac{\rho^r}{(100 - r)}, \quad \pi_w[90] = W \ln 10, \quad \pi_w[95] = W \ln 20$$

$$W_q[t] = P[q \leq t] = 1 - \rho \exp\left(\frac{-t}{W}\right), \quad P[q > t] = \rho \exp\left(\frac{-t}{W}\right)$$

$$W_q = \frac{\rho W_s}{1 - \rho}, \quad \sigma_q^2 = \frac{(2 - \rho)\rho W_s^2}{(1 - \rho)^2}$$

$$\pi_q[r] = \max\left\{W \ln\left(\frac{100\rho}{100 - r}\right), 0\right\}$$

$$\pi_q[90] = \max\{W \ln(10\rho), 0\}, \quad \pi_q[95] = \max\{W \ln(20\rho), 0\}$$

Table 4. M/M/1/K Queueing System

$$p_n = \begin{cases} \frac{(1 - a)a^n}{(1 - a^{K+1})} & \text{if } \lambda \neq \mu, \\ \frac{1}{K + 1} & \text{if } \lambda = \mu, \end{cases}$$

$n = 0, 1, \dots, K$, where $a = \lambda W_s$.

$\lambda_a = (1 - p_K)\lambda$, Mean arrival rate into system.

$$L = \begin{cases} \frac{a[1 - (K + 1)a^K + Ka^{K+1}]}{(1 - a)(1 - a^{K+1})} & \text{if } \lambda \neq \mu, \\ L = \frac{\rho}{1 - \rho} & \text{if } \lambda = \mu. \end{cases}$$

$$L_q = L - (1 - p_0), \quad q_n = \frac{p_n}{1 - p_K}, \quad n = 0, 1, \dots, K - 1.$$

$$W[t] = 1 - \sum_{n=0}^{K-1} q_n Q[n; \mu t],$$

where

$$q[n; \mu t] = e^{-\mu t} \sum_{k=0}^n \frac{\mu t^k}{k!}$$

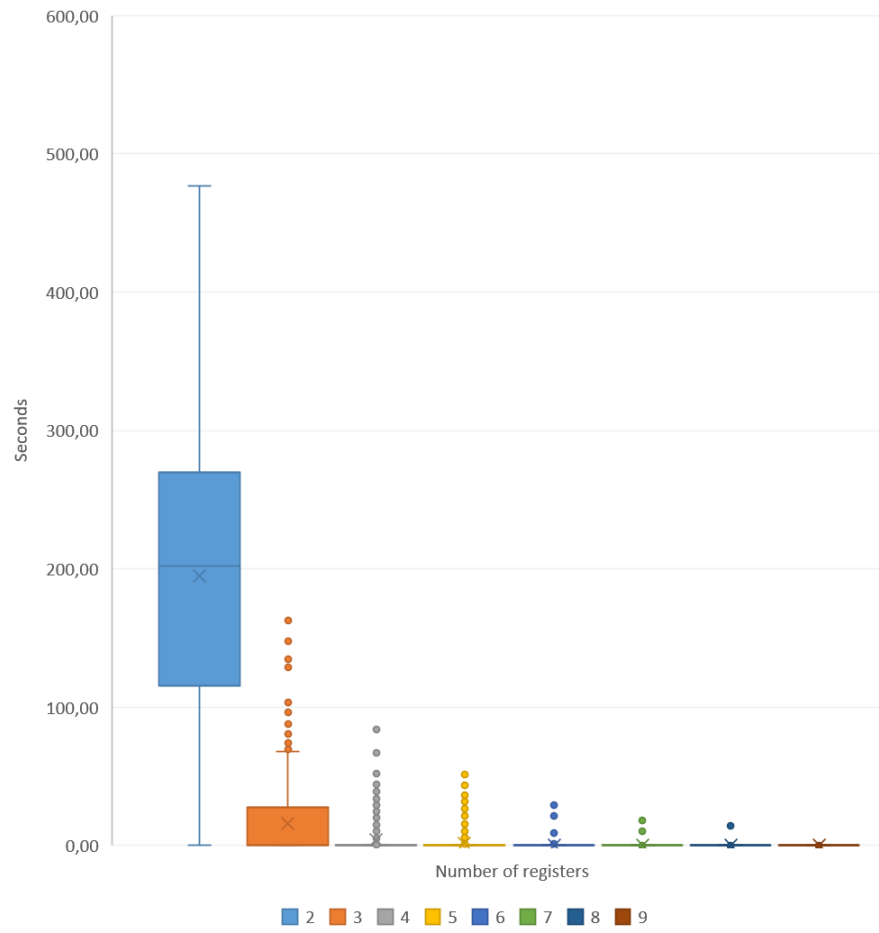
$$W = \frac{L}{\lambda_a}, \quad W_q = \frac{L_q}{\lambda_a}$$

$$W_q[t] = 1 - \sum_{n=0}^{K-2} q_{n+1} Q[n; \mu t].$$

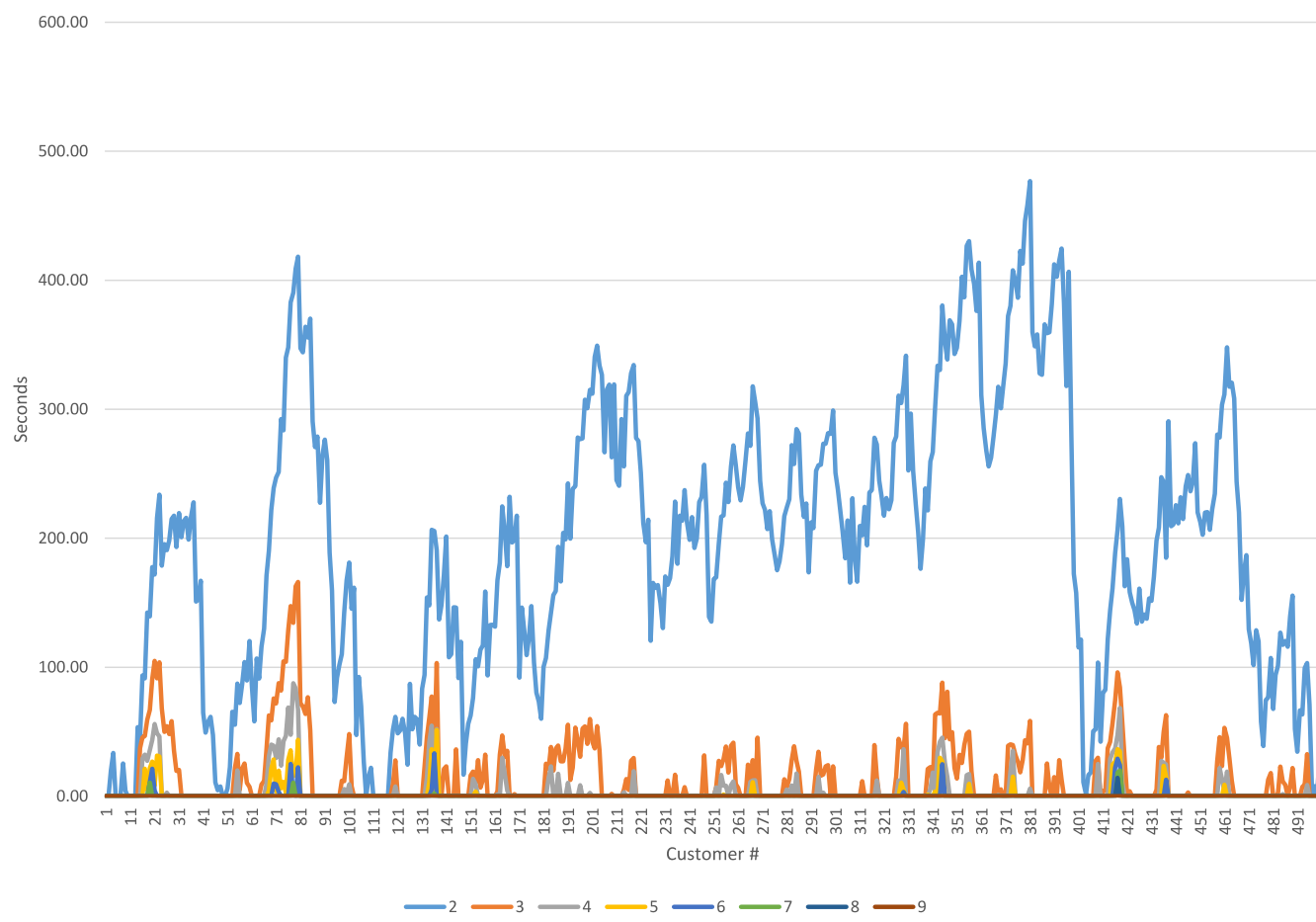
$$E[q | q > 0] = \frac{W_q}{1 - p_0}, \quad \rho = (1 - p_K)a.$$

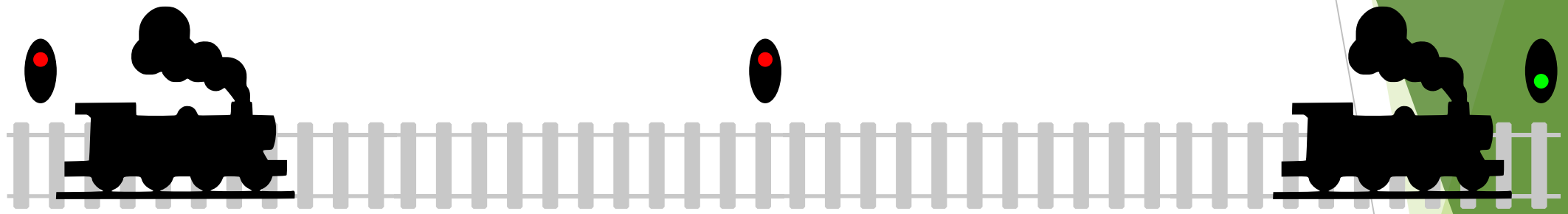
16 more types of queueing systems

Queueing times

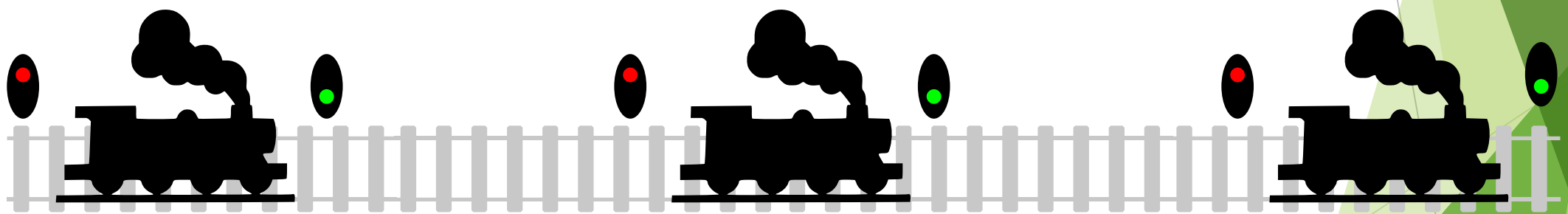


Queueing times

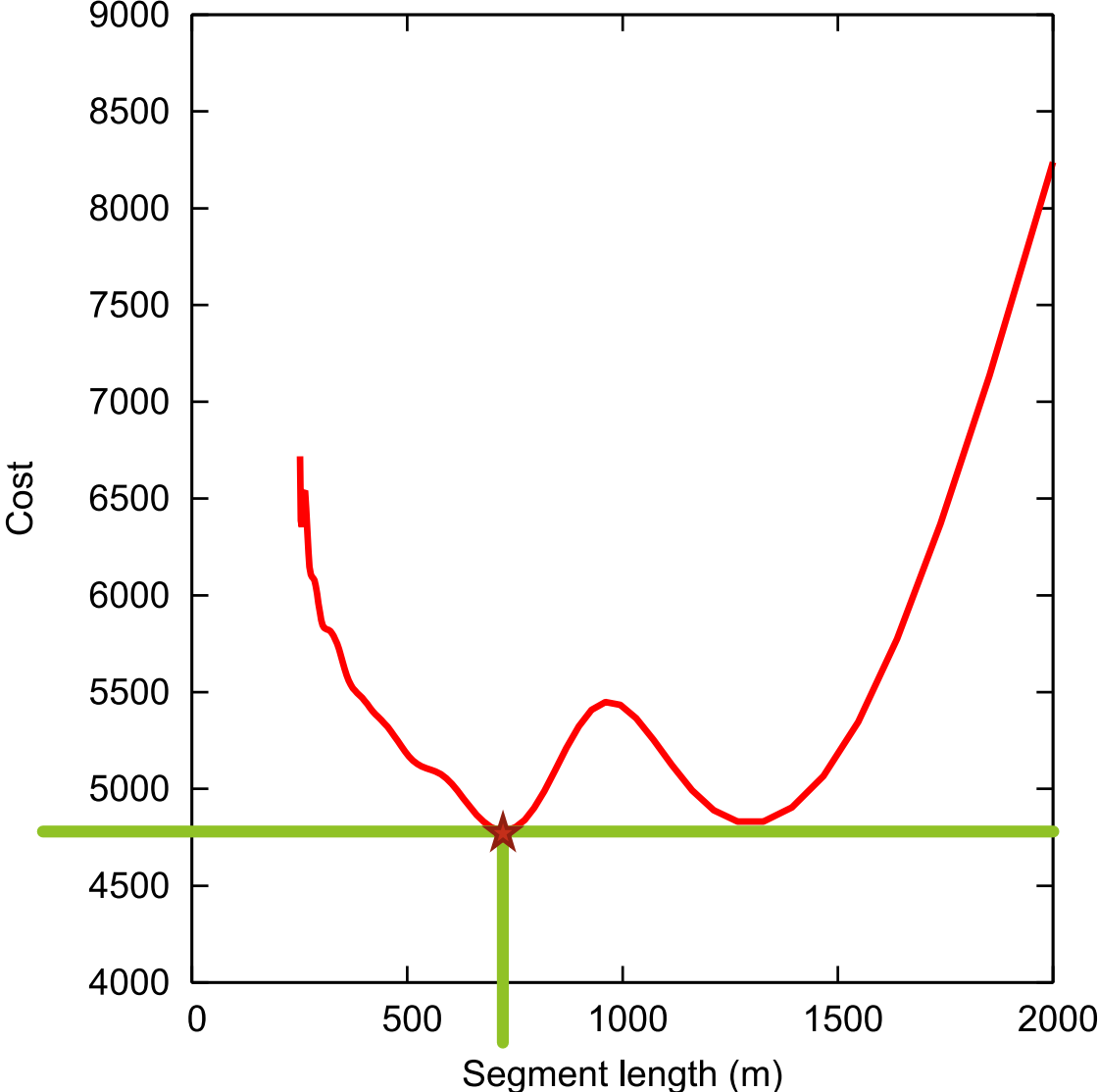


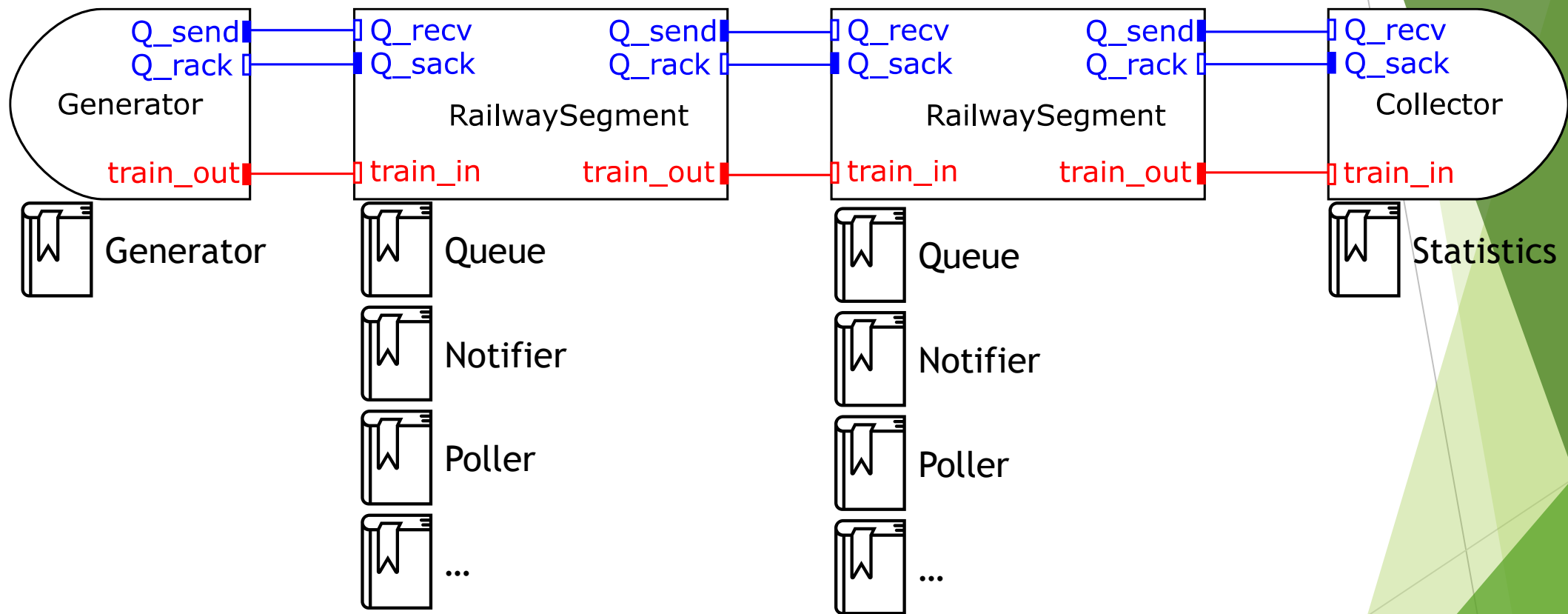


VS.



$$cost = 10 \times \#lights + avg(t_{travel})$$





Custom Atomic DEVS Models

Extending a DEVS repository

Atomic DEVS Specification

$$M = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

X : set of input events

Y : set of output events

S : set of sequential states

$q_{init} : Q$

$$Q = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$$

$\delta_{int} : S \rightarrow S$

$\delta_{ext} : Q \times X \rightarrow S$

$\lambda : S \rightarrow Y \cup \{\phi\}$

$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$

Atomic DEVS Specification

$$M = \langle X, Y, S, q_{init}, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

X : set of input events

Y : set of output events

S : set of states

PATTERNS!

$$P = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$$

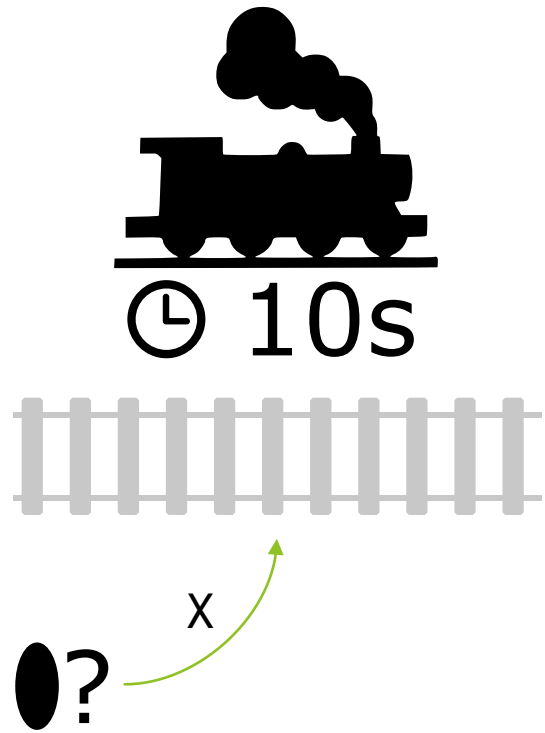
$$\delta_{int} : S \rightarrow S$$

$$\delta_{ext} : Q \times X \rightarrow S$$

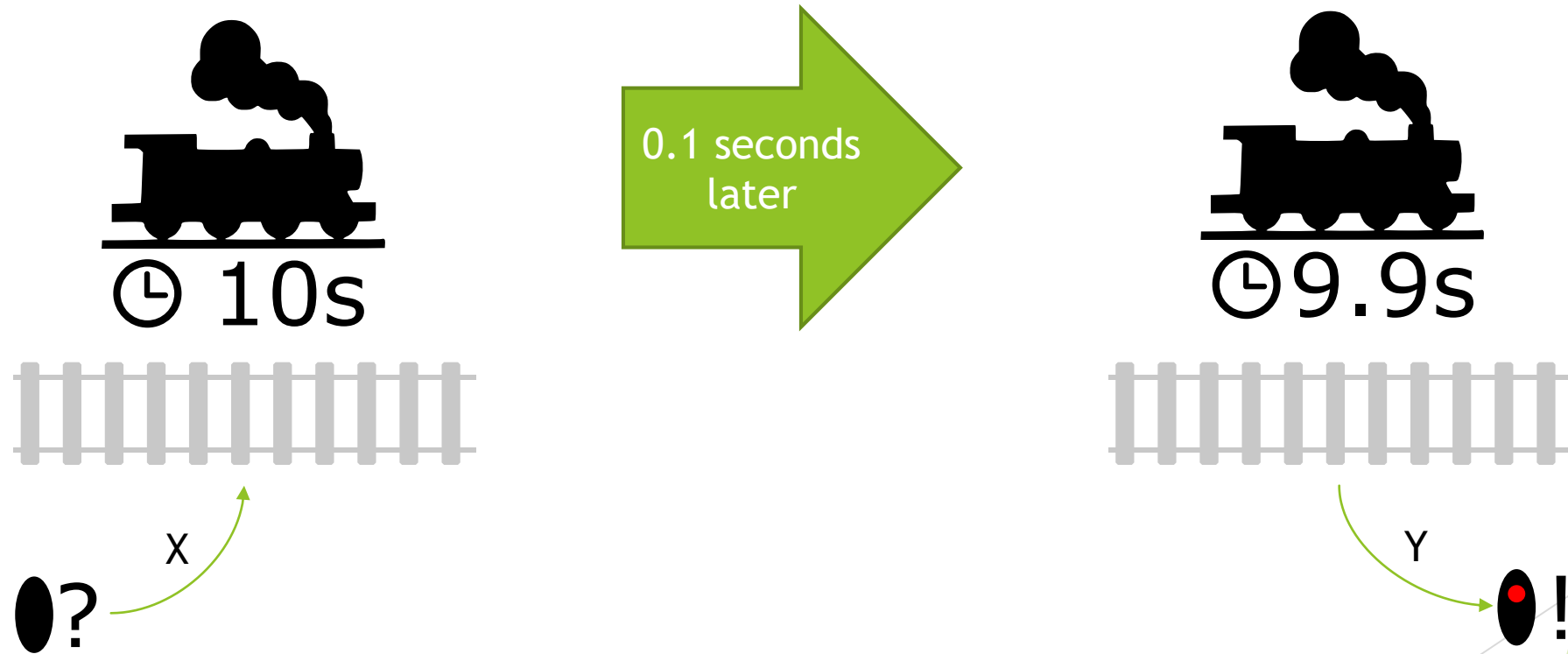
$$\lambda : S \rightarrow Y \cup \{\phi\}$$

$$ta : S \rightarrow \mathbb{R}_{0,+\infty}^+$$

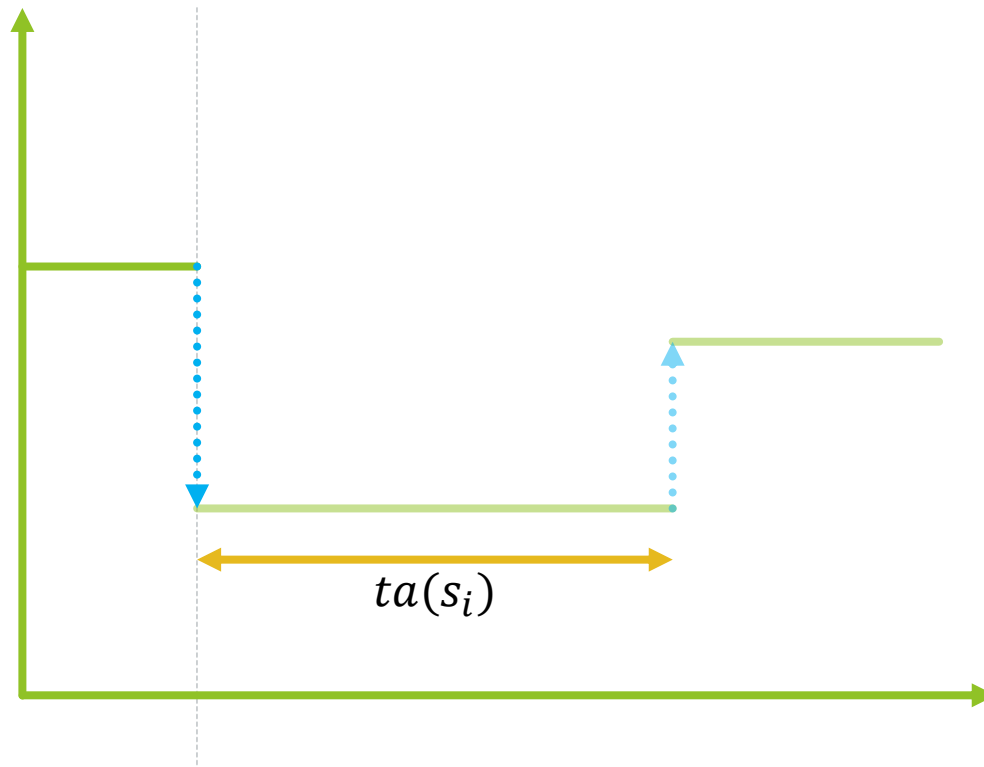
Pattern 1: Ignore an Event



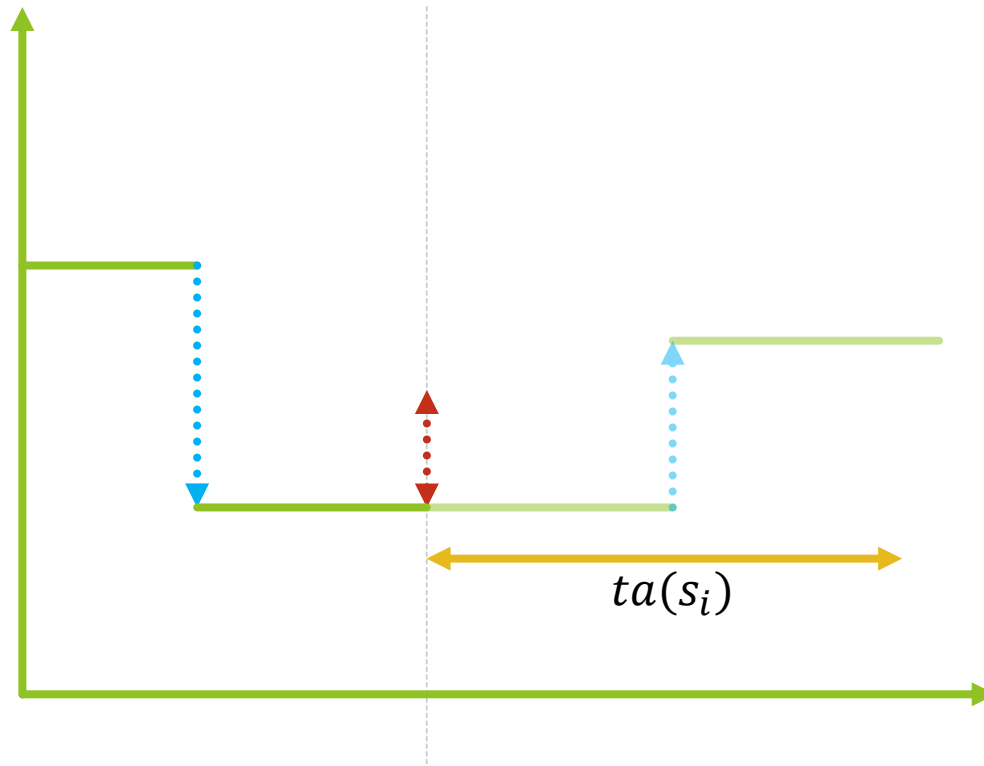
Pattern 1: Ignore an Event



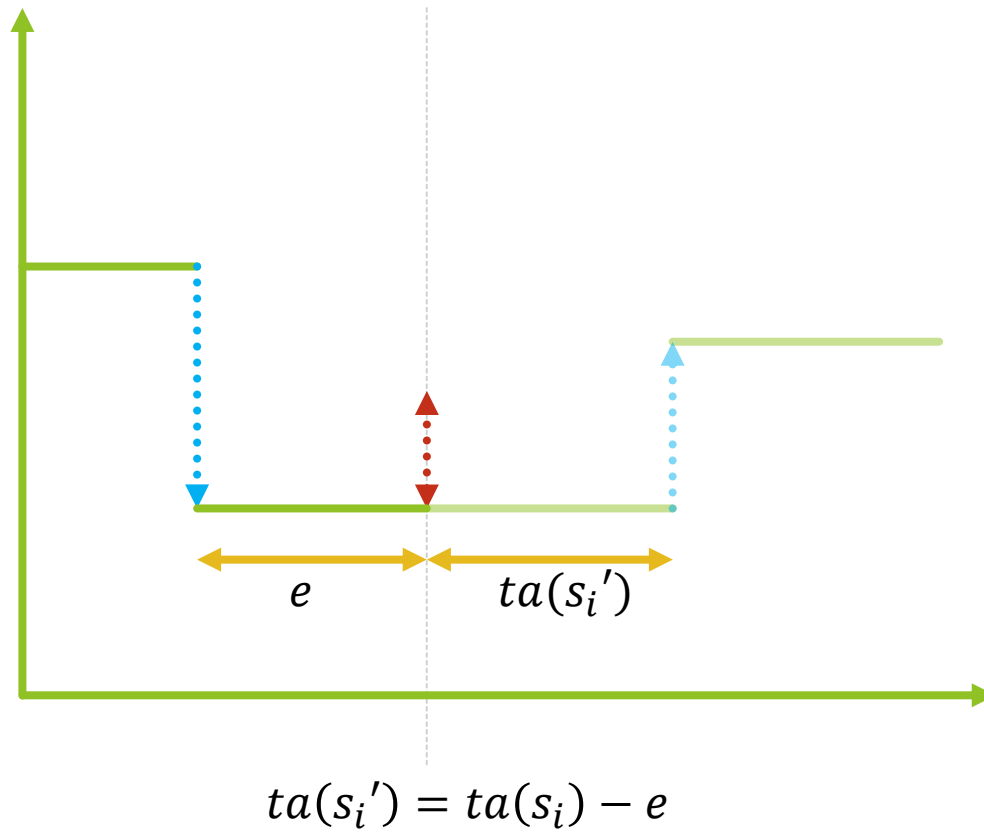
Pattern 1: Ignore an Event



Pattern 1: Ignore an Event



Pattern 1: Ignore an Event



Pattern 1: Ignore an Event

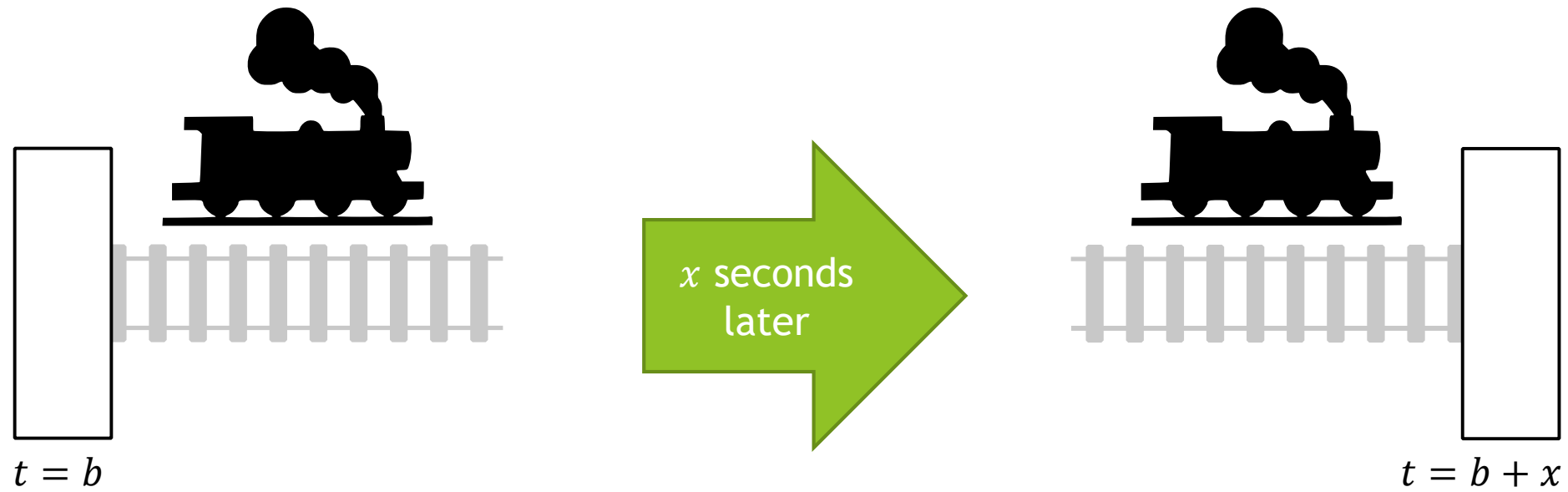


```
class RailwaySegmentState():
    def __init__(self):
        self.timer = INFINITY

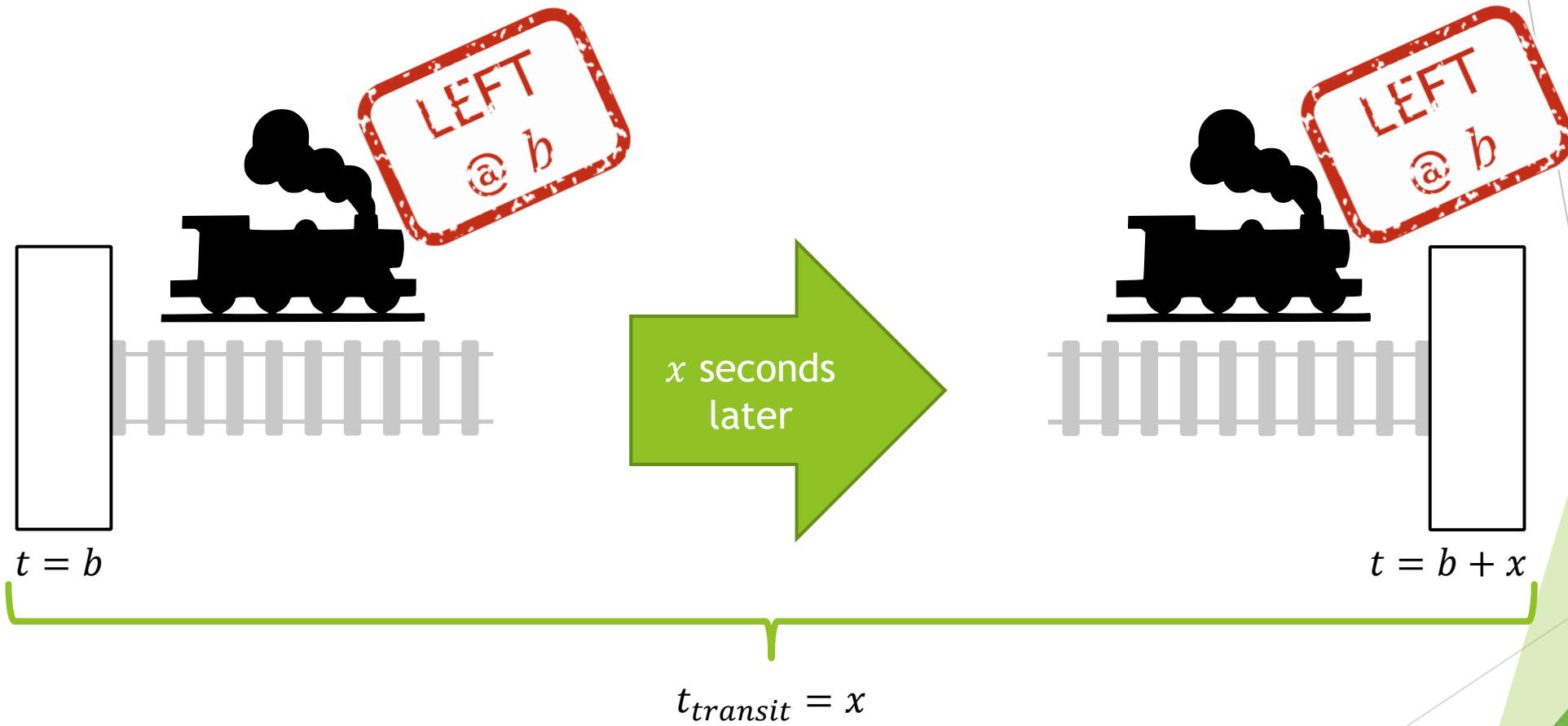
class RailwaySegment(AtomicDEVS):
    def extTransition(self, inputs):
        self.state.timer -= self.elapsed
        ...

    def timeAdvance(self):
        return self.state.timer
```

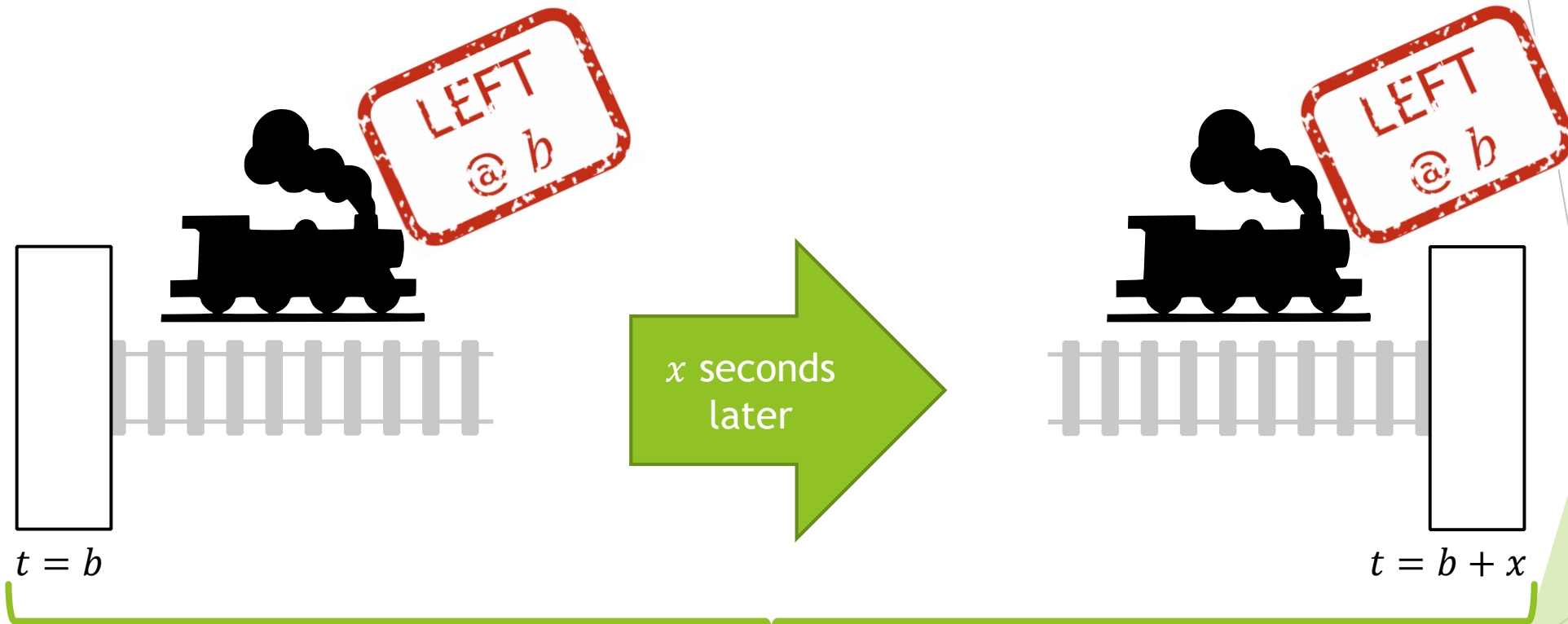
Pattern 2: Simulated Time



Pattern 2: Simulated Time



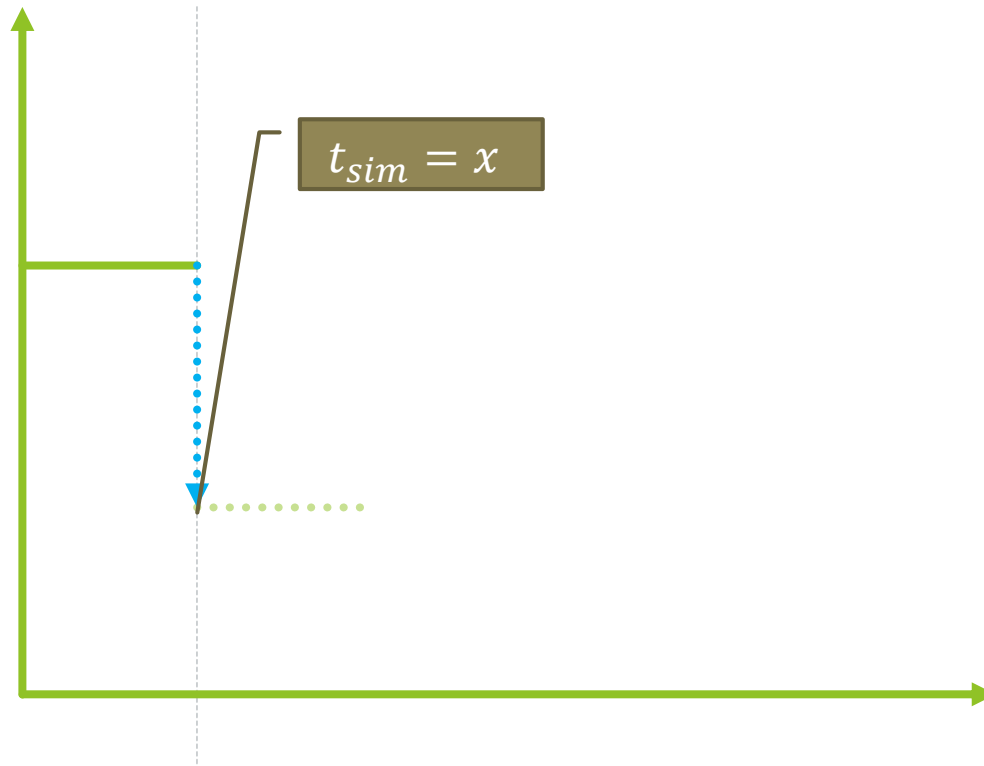
Pattern 2: Simulated Time



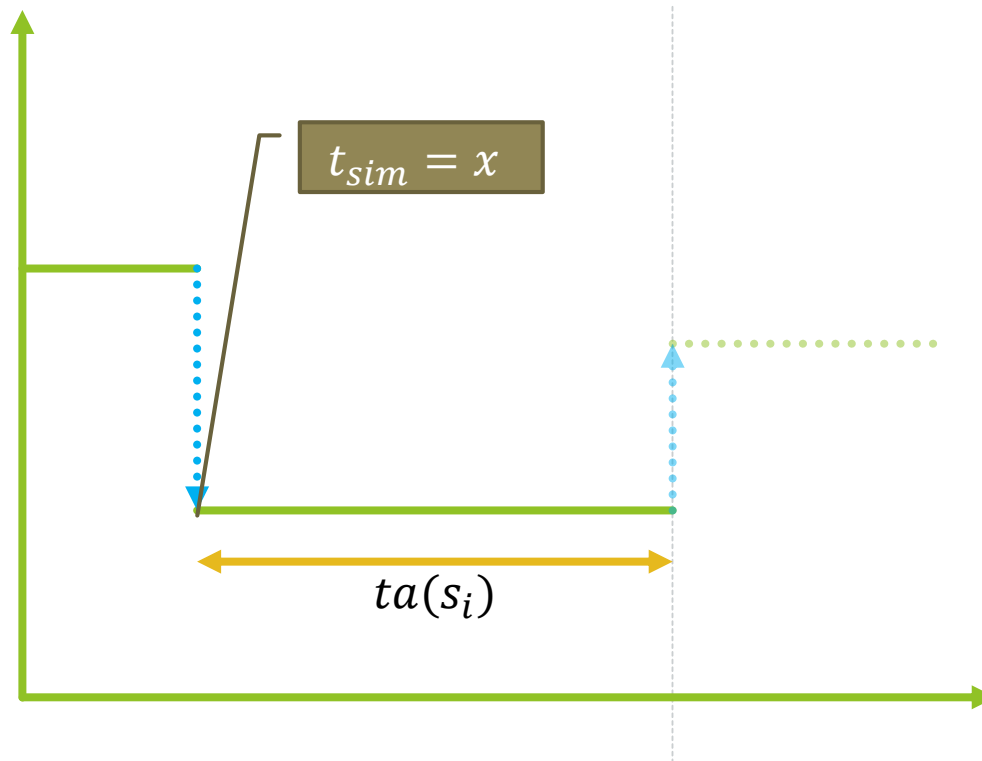
$$t_{transit} = x$$

$$t = ?$$

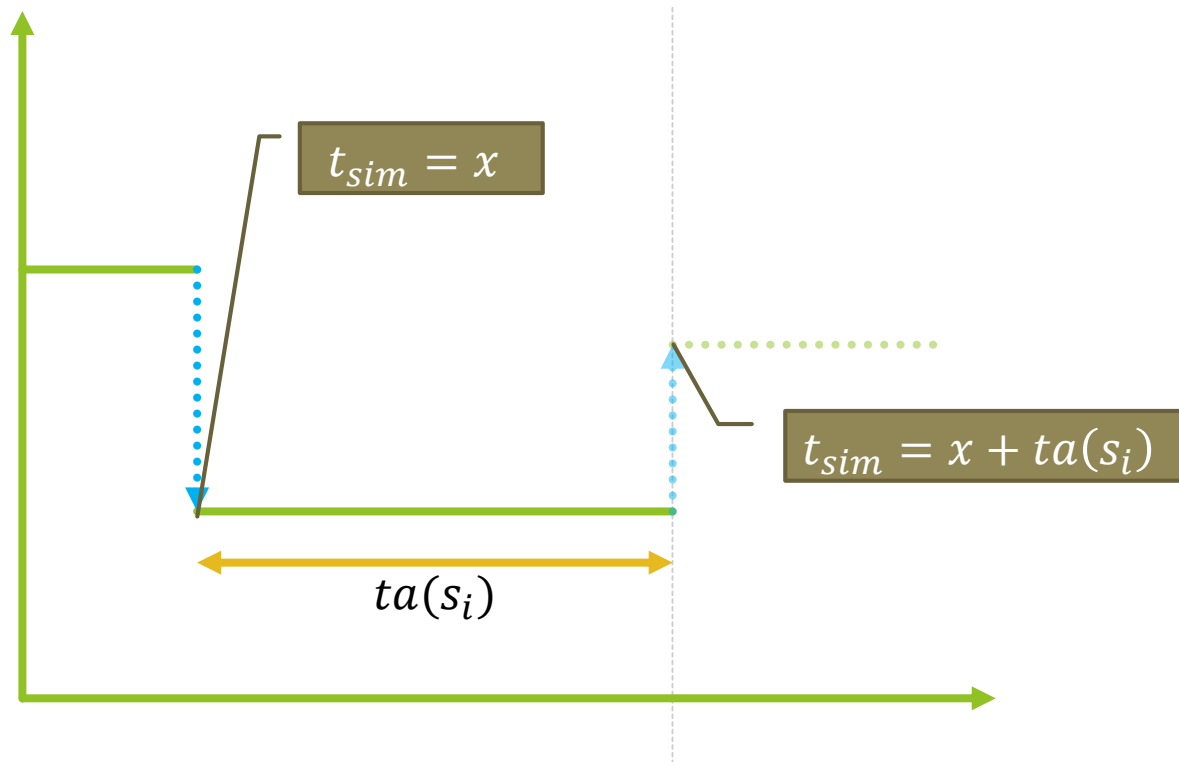
Pattern 2: Simulated Time



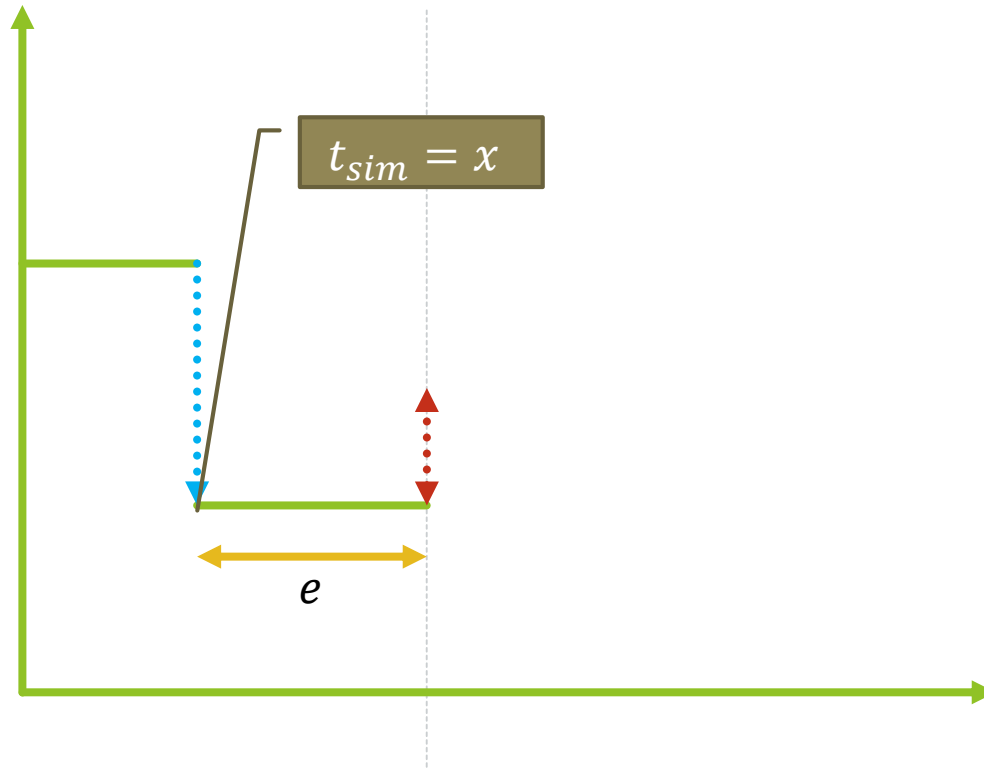
Pattern 2: Simulated Time



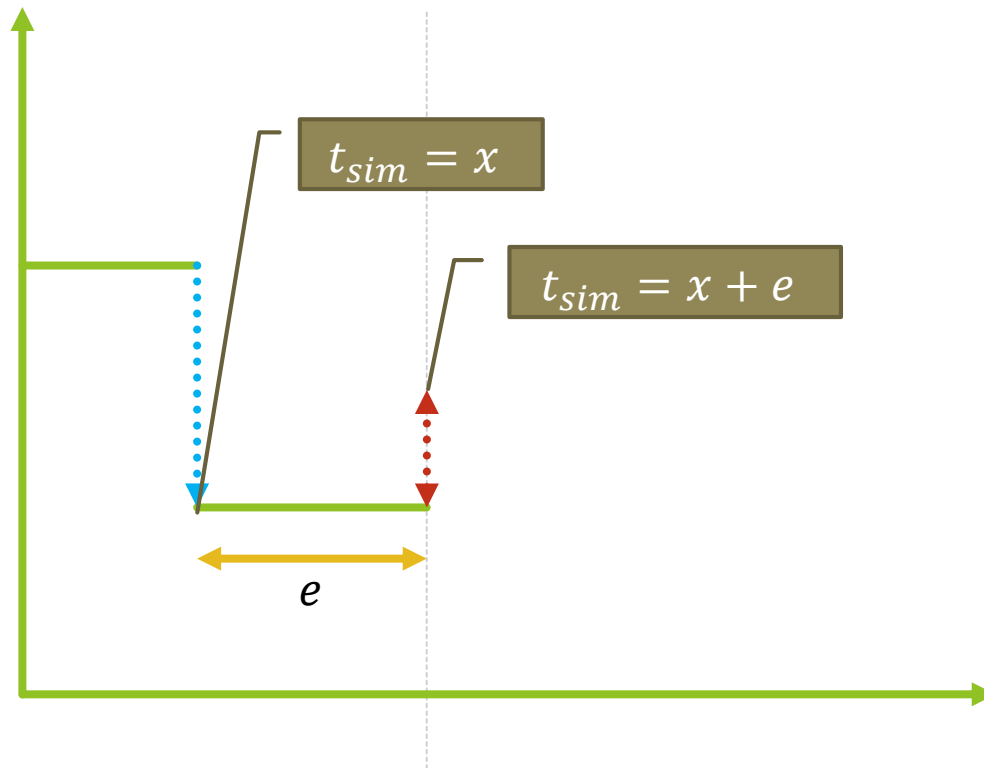
Pattern 2: Simulated Time



Pattern 2: Simulated Time



Pattern 2: Simulated Time



Pattern 2: Simulated Time



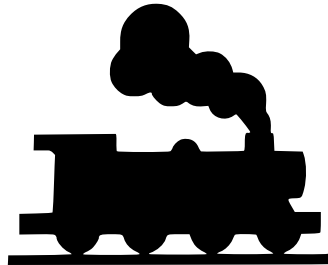
```
class GeneratorState:
    def __init__(self):
        self.t_sim = 0.0
        ...

class Generator(AtomicDEVS):
    def __init__(self):
        self.state = GeneratorState()
        ...

    def intTransition(self):
        self.state.t_sim += self.timeAdvance()
        ...

    def extTransition(self, inputs):
        self.state.t_sim += self.elapsed
        ...
```

Pattern 3: Multiple Timers

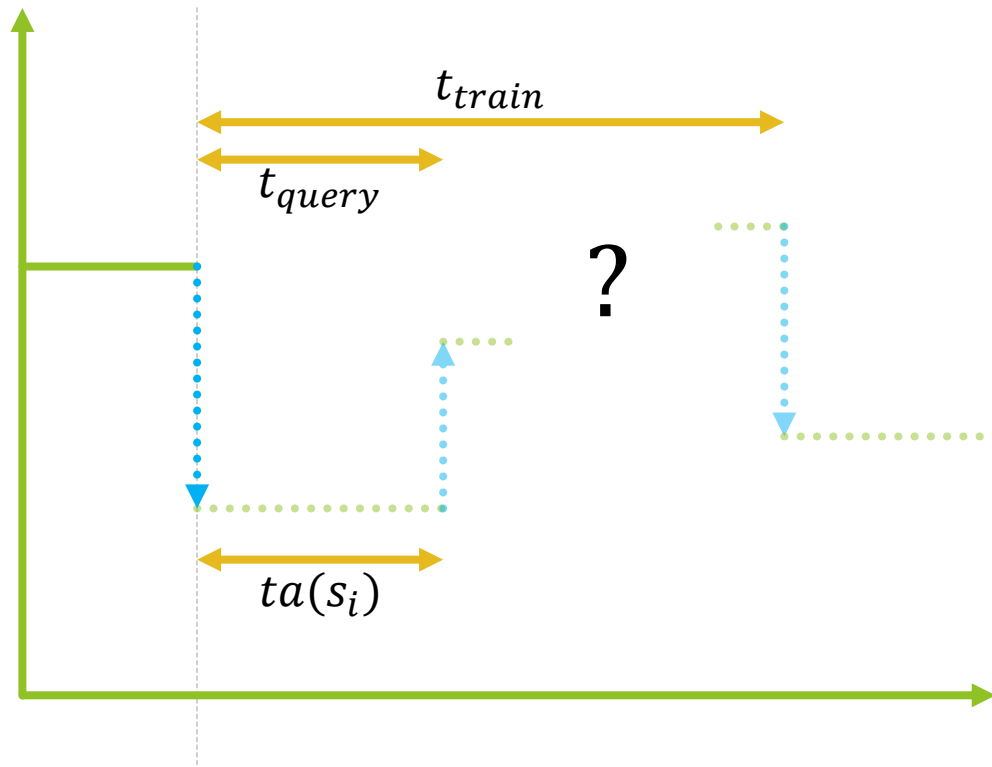


🕒 10s

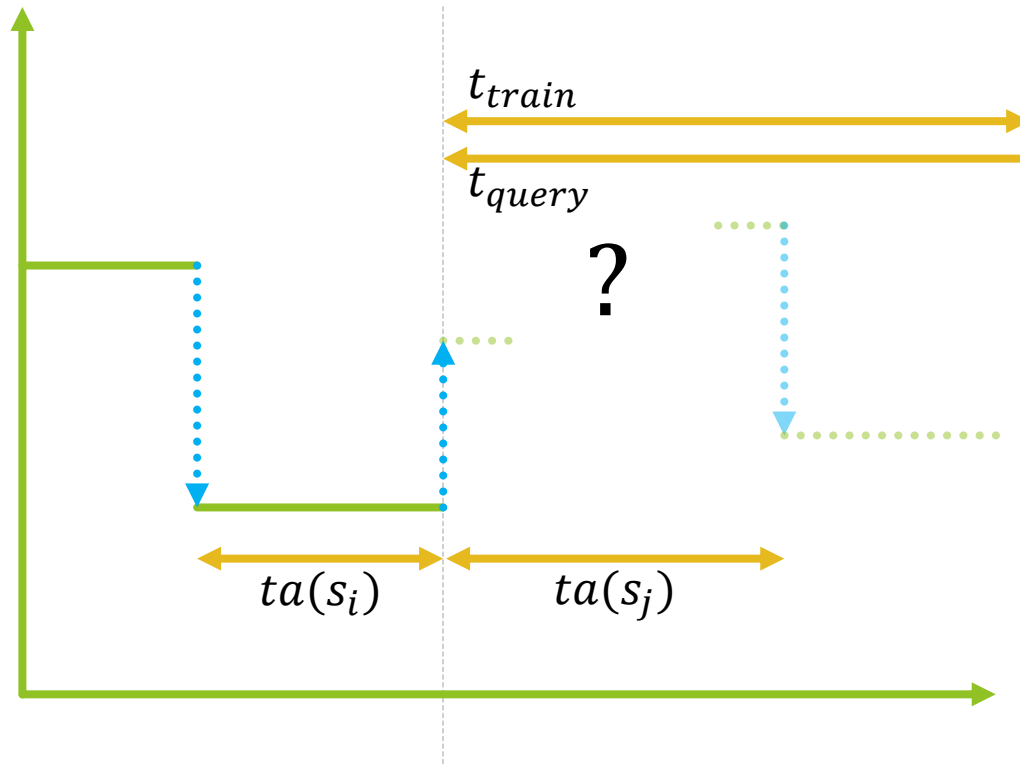
●?🕒 0.1s



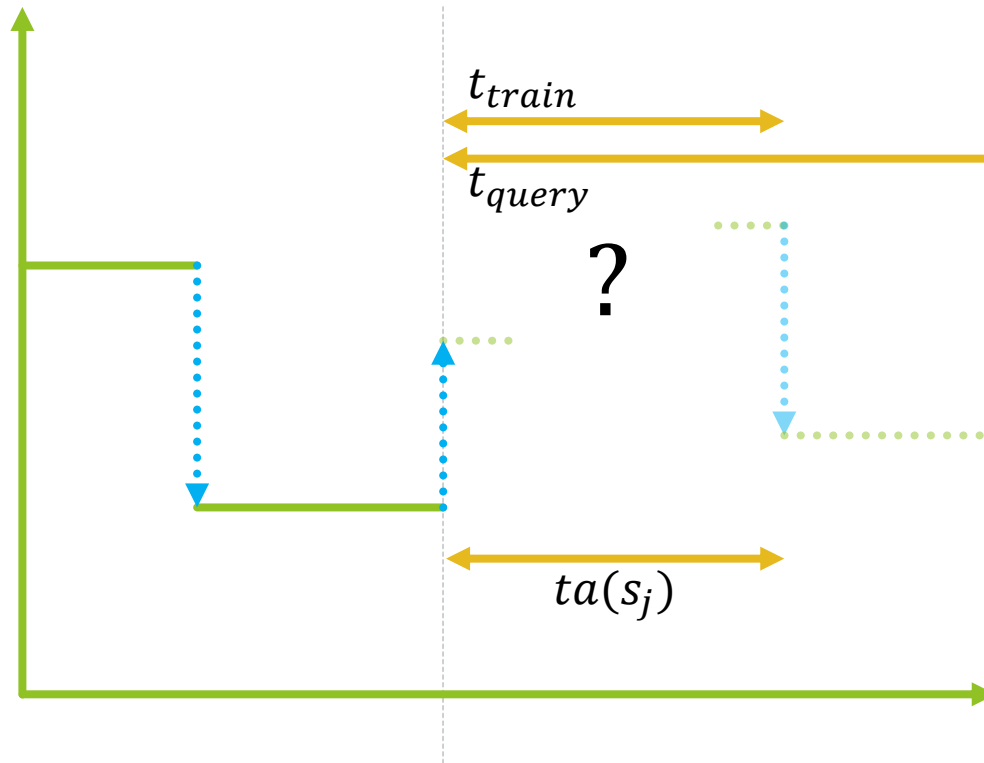
Pattern 3: Multiple Timers



Pattern 3: Multiple Timers



Pattern 3: Multiple Timers



Pattern 3: Multiple Timers



```
class RailwaySegmentState:
    def __init__(self):
        self.t_query = INFINITY
        self.t_train = INFINITY

class RailwaySegment(AtomicDEVS):
    def timeAdvance(self):
        return min(self.state.t_query, \
                  self.state.t_train)

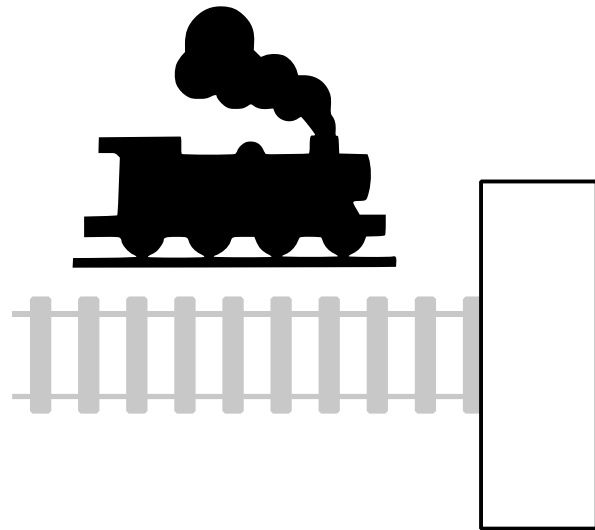
    def extTransition(self, inputs):
        self.state.t_query -= self.elapsed
        self.state.t_train -= self.elapsed
    ...
```



```
...

def intTransition(self):
    self.state.t_query -= self.timeAdvance()
    self.state.t_train -= self.timeAdvance()
    if (self.state.t_query == 0):
        ... # process query
    elif (self.state.t_train == 0):
        ... # process train
```

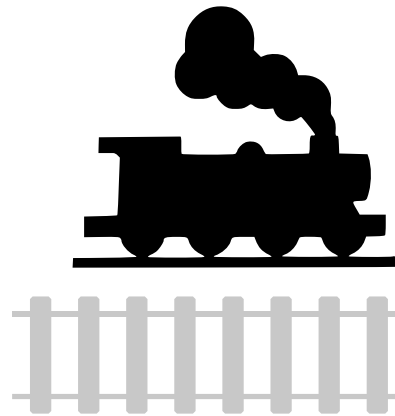

Pattern 4: Statistics Gathering



n trains = n doubles = $8n$ bytes

[1815.7; 1901.1; 1801.4; 1867.3; 1911.8; 1846.4; 1844.3; 1873.5; ...]

Pattern 4: Statistics Gathering



n trains = n doubles = $8n$ bytes

[1815.7; 1901.1; 1801.4; 1867.3; ~~181.8~~; 1846.4; 1844.3; 1873.5; ...]



$$sum = \sum_i t_{travel} = 36,902,140.45$$

$$count = |arrivals| = 20,000$$

$$avg = \frac{sum}{count}$$

n trains = 1 double + 1 long long int = 16 bytes

Pattern 4: Statistics Gathering

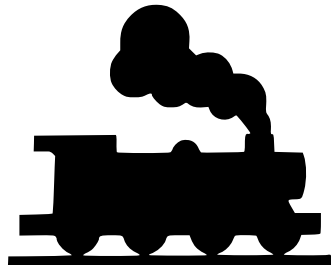


```
class CollectorState():
    def __init__(self):
        self.counter = 0
        self.accum = 0.0
        self.t_sim = 0.0

class Collector(AtomicDEVS):
    def extTransition(self, inputs):
        left = inputs[self.train].left
        transit = self.state.t_sim - left
        self.accum += transit
        self.counter += 1
```

Pattern 5: Complex State/Event

S

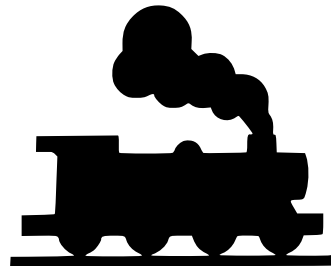


⌚ 10s

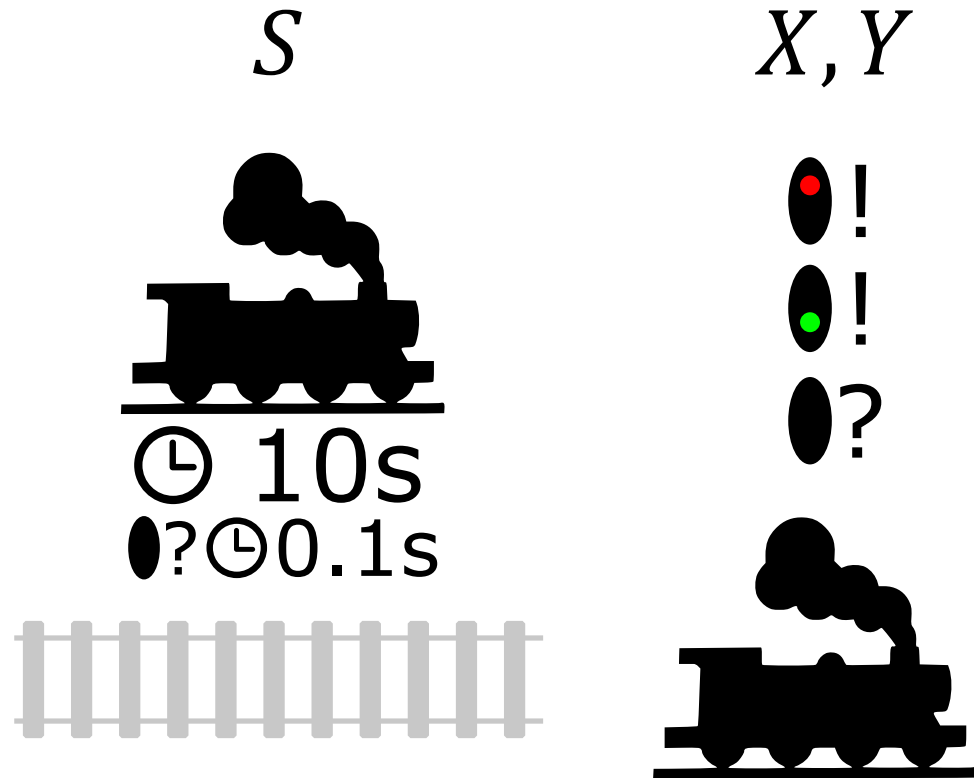
●? ⌚ 0.1s



X, Y



Pattern 5: Complex State/Event



$S = \{ \langle timer_{train}, timer_{query}, v_{train}, t_{train}, a_{train} \rangle \mid timer_{train} \in \mathbb{R}^+, \dots \}$

$X, Y = \{ \langle t, v, a \rangle \mid t \in \mathbb{R}^+, \dots \} \cup \{ query \} \cup \{ colour \mid colour \in \{ red, green \} \}$

Pattern 5: Complex State/Event



```
class RailwaySegmentState:
    def __init__(self):
        self.train = None

class RailwaySegment(AtomicDEVS):
    def __init__(self):
        self.state = RailwaySegmentState()
        ...

    def extTransition(self, inputs):
        ...
        self.state.train = inputs[self.train_in]
        ...
```



```
class Train:
    def __init__(self, t, a):
        self.t = t
        self.a = a
        self.v = 0.0

class Query:
    def __init__(self):
        pass

class QueryAck:
    def __init__(self, colour):
        self.colour = colour
```

Conclusion

